

Gigaflow: Pipeline-Aware Sub-Traversal Caching for Modern SmartNICs

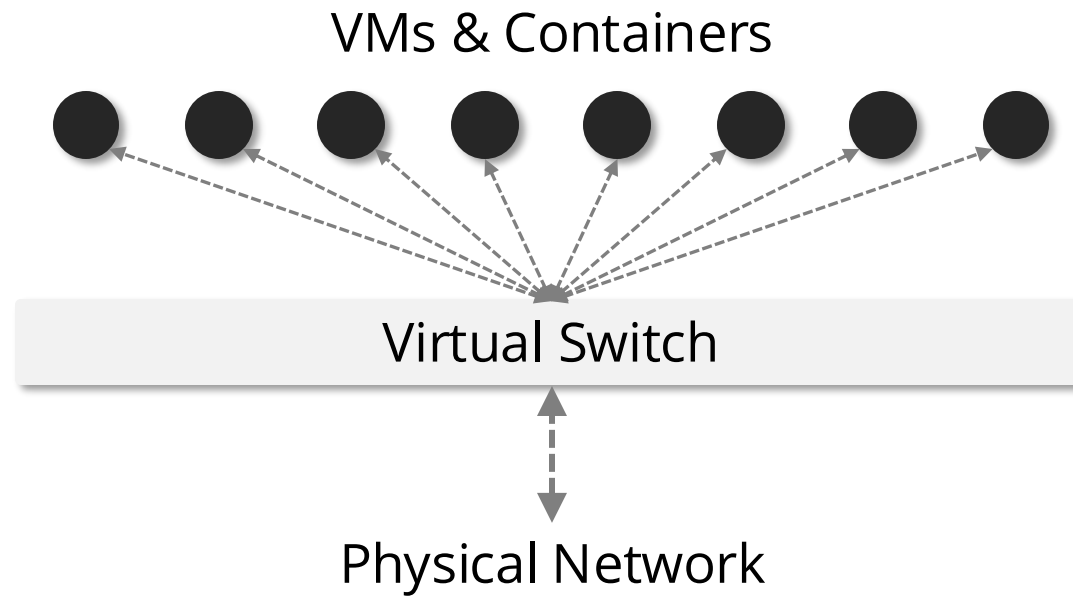
Annus Zulfiqar

Ali Imran, Venkat Kunaparaju, Ben Pfaff,
Gianni Antichi, Muhammad Shahbaz

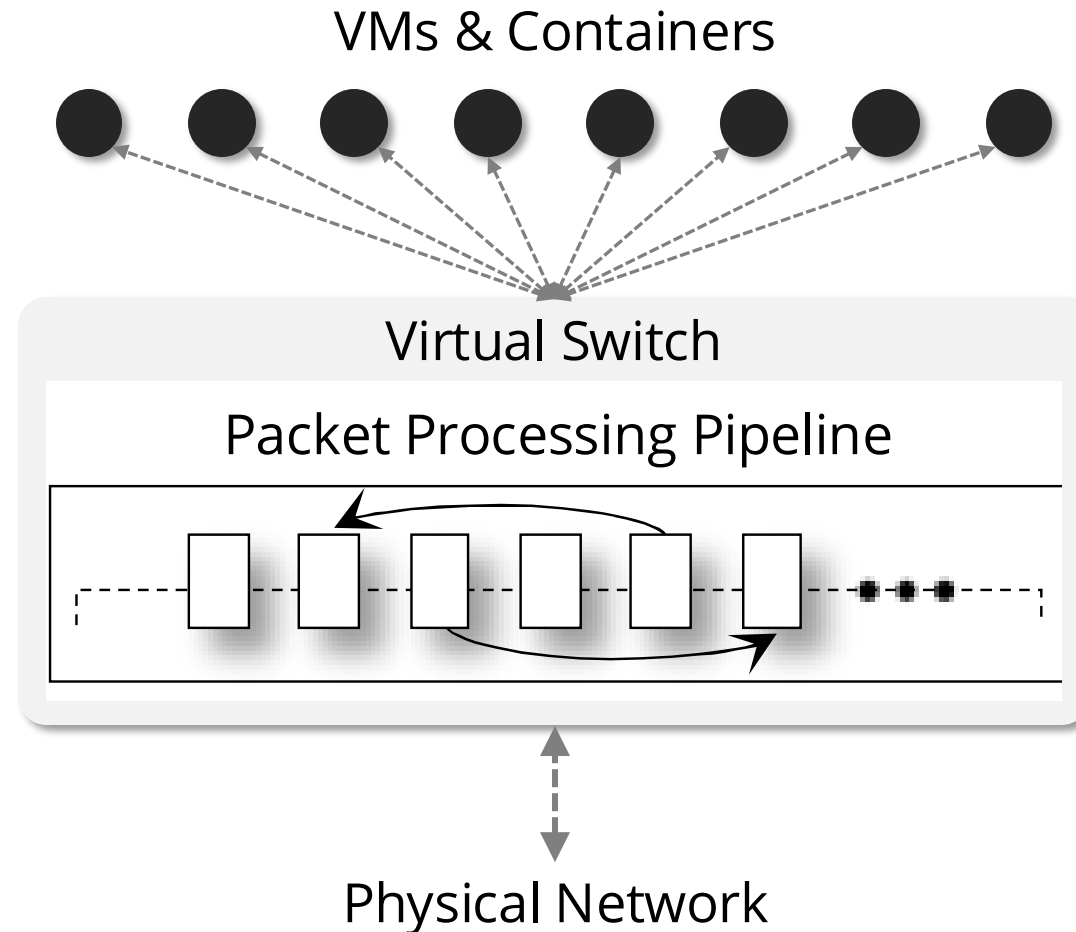


POLITECNICO
MILANO 1863

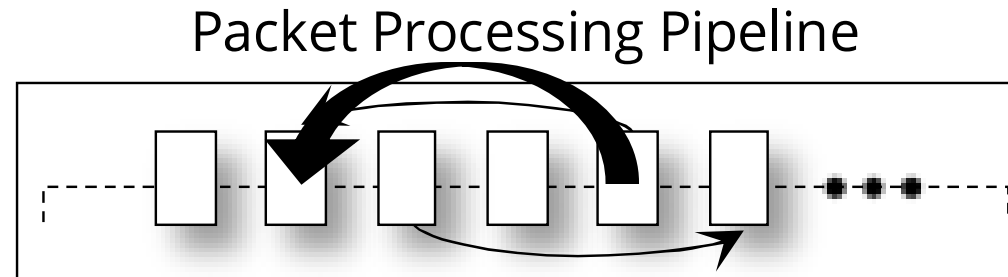
Virtual Switches in the Data Center



Virtual Switches in the Data Center

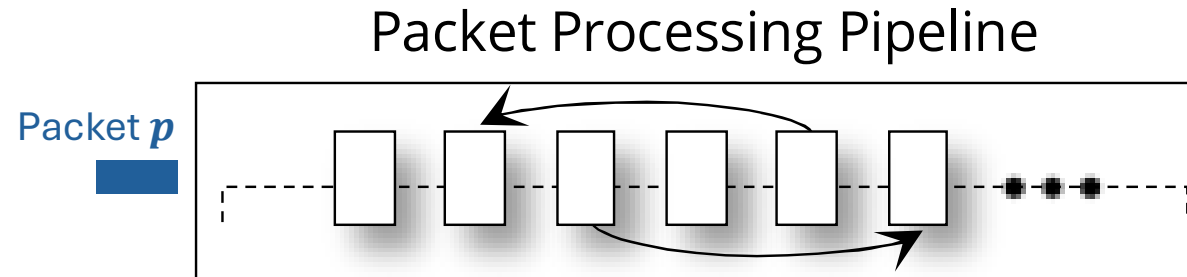


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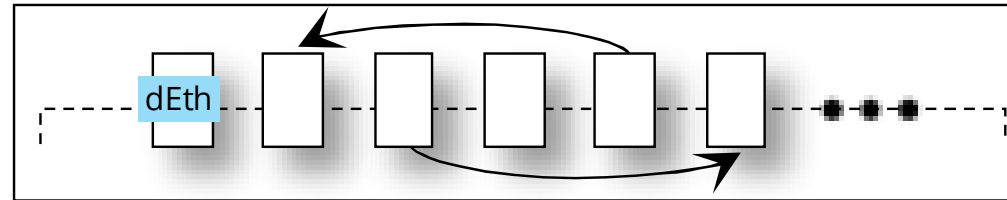
Some networking policies require **recursive application of rules**, e.g., encapsulation and decapsulation actions

Virtual Switches in the Data Center

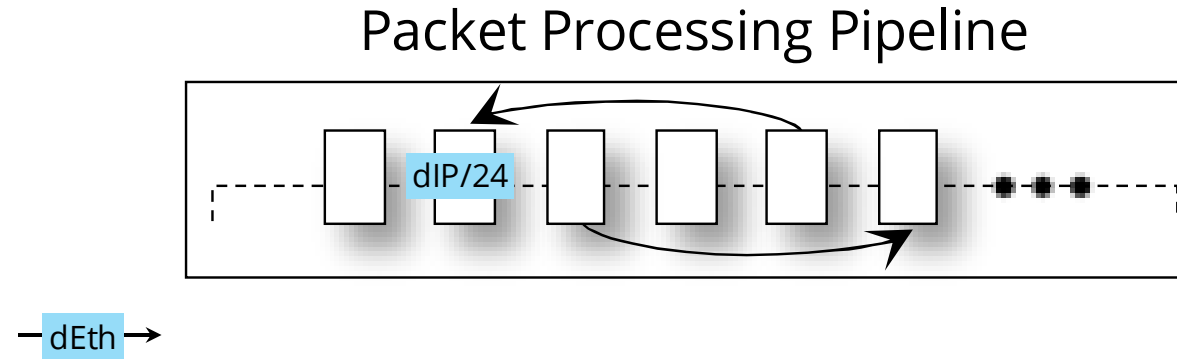


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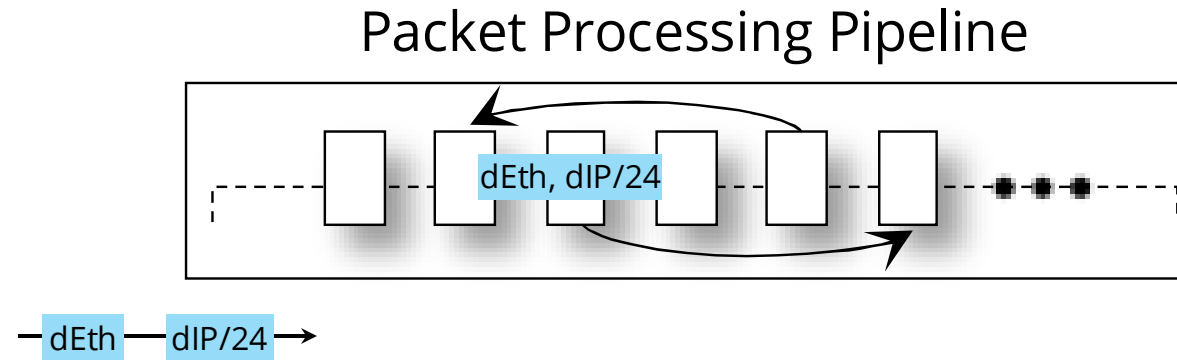
Packet Processing Pipeline



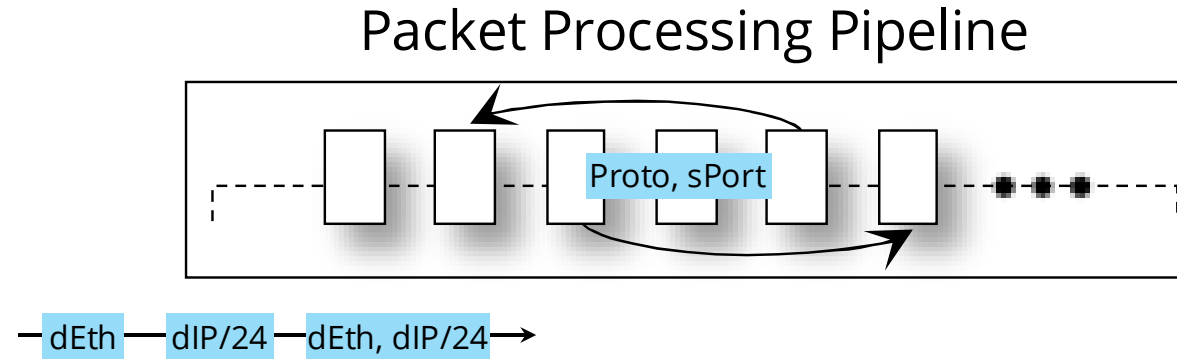
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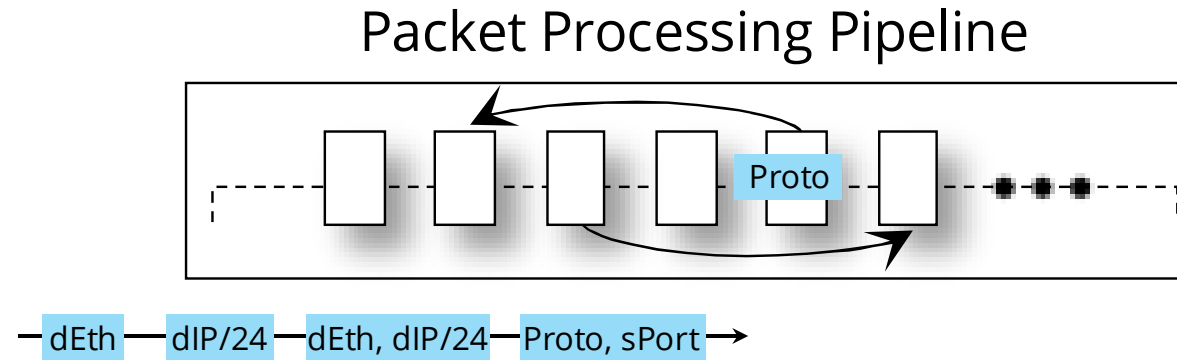
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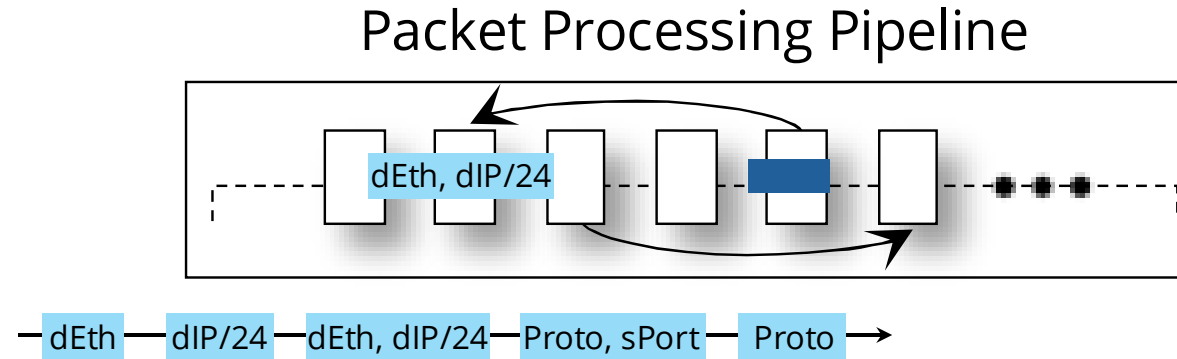
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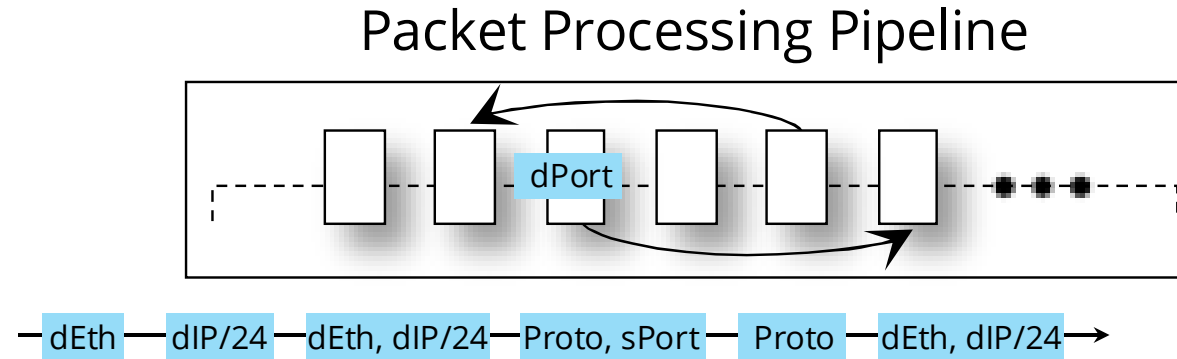
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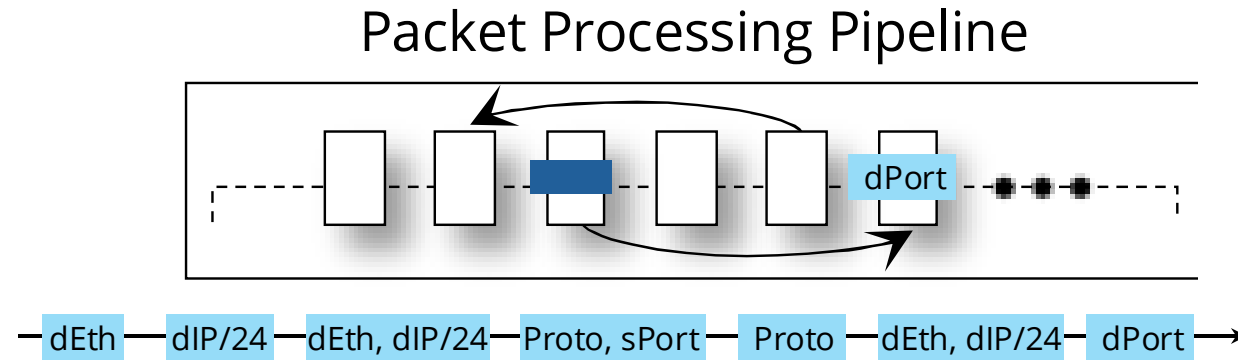
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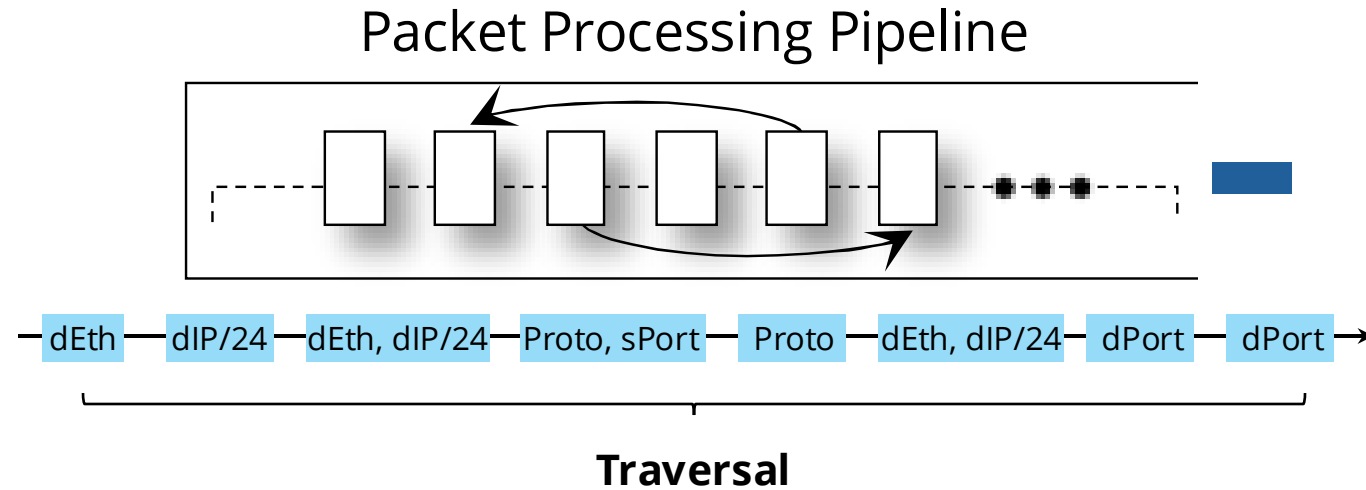
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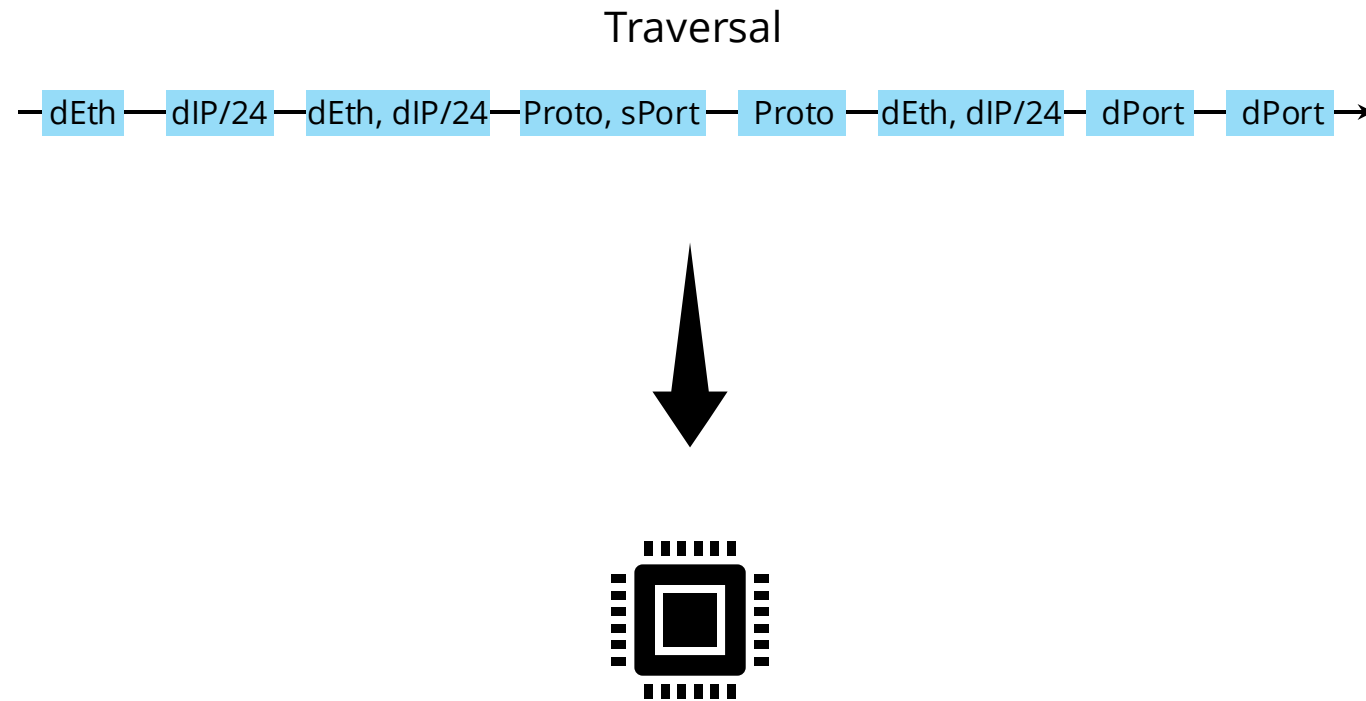
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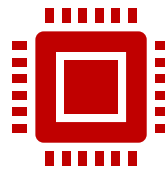


Virtual Switches in the Data Center



Virtual Switches in the Data Center

Multi-table lookups are expensive

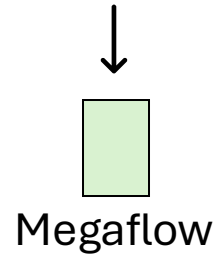


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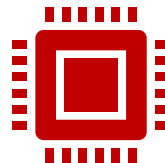
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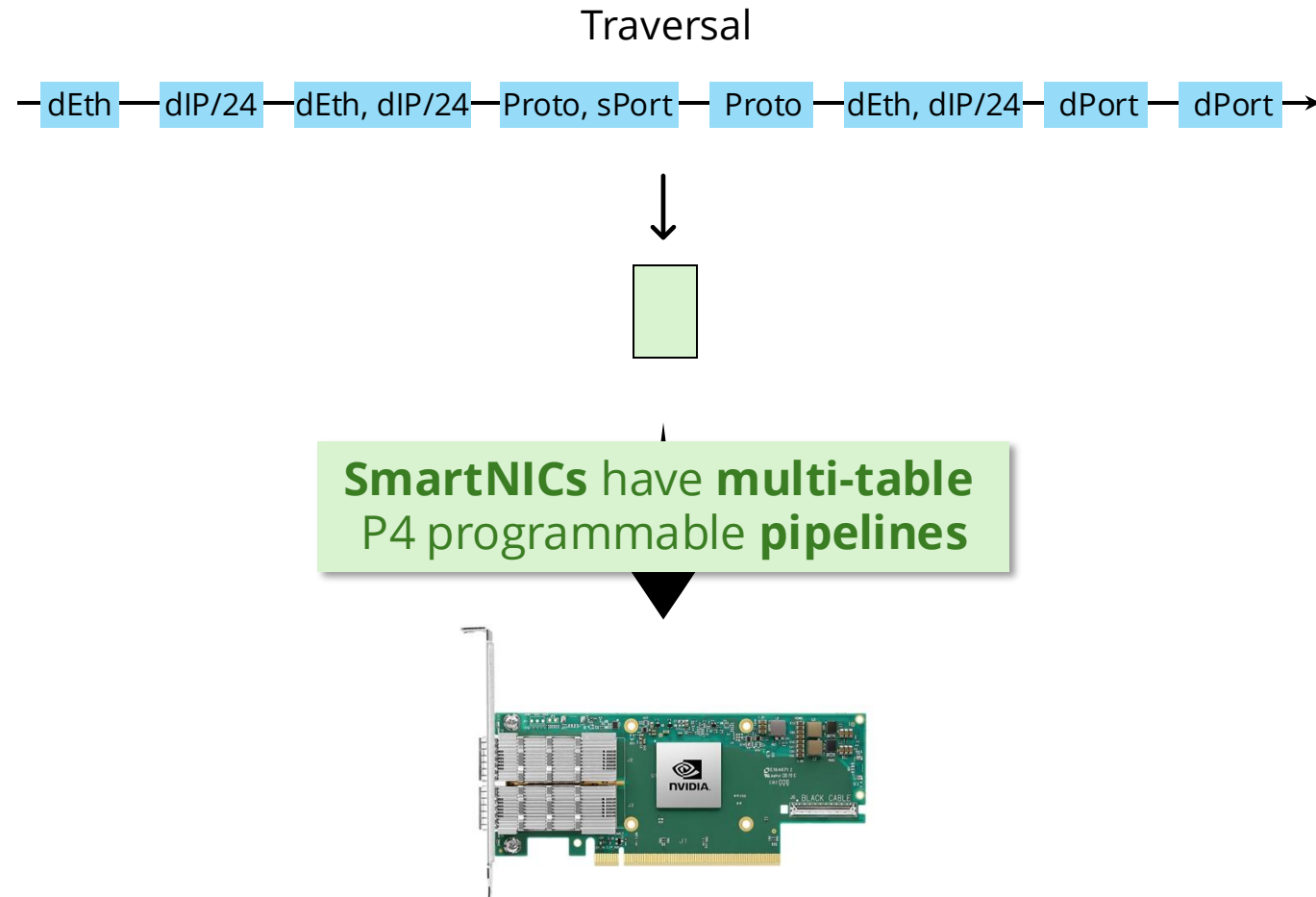
Caches must be single table
for lookup speed



The traffic alone will guide
cache **rule generation**



The Target Architecture has Changed

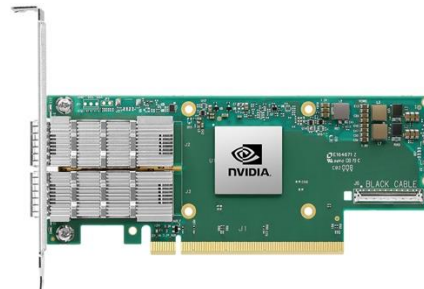


The Target Architecture has Changed

Traversal

—dEth—dIP/24—dEth, dIP/24—Proto, sPort—Proto—dEth, dIP/24—dPort—dPort→

It's time for a **clean slate**
virtual switch design



Offload the vSwitch to SmartNIC Tables?

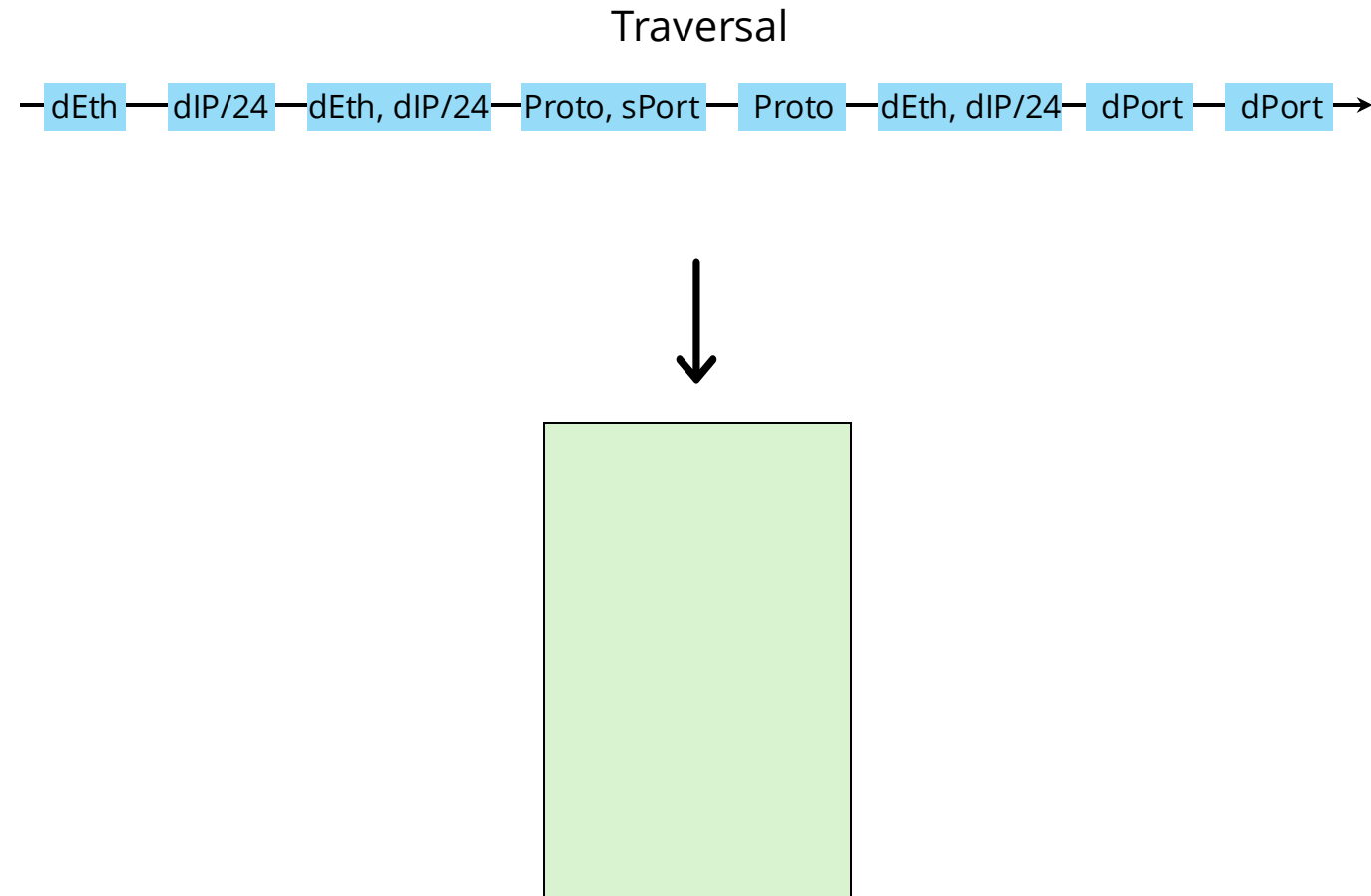
vSwitches provide an
infinite resource abstraction
to the controller

**We can't limit the vSwitch
packet pipeline's flexibility!!**

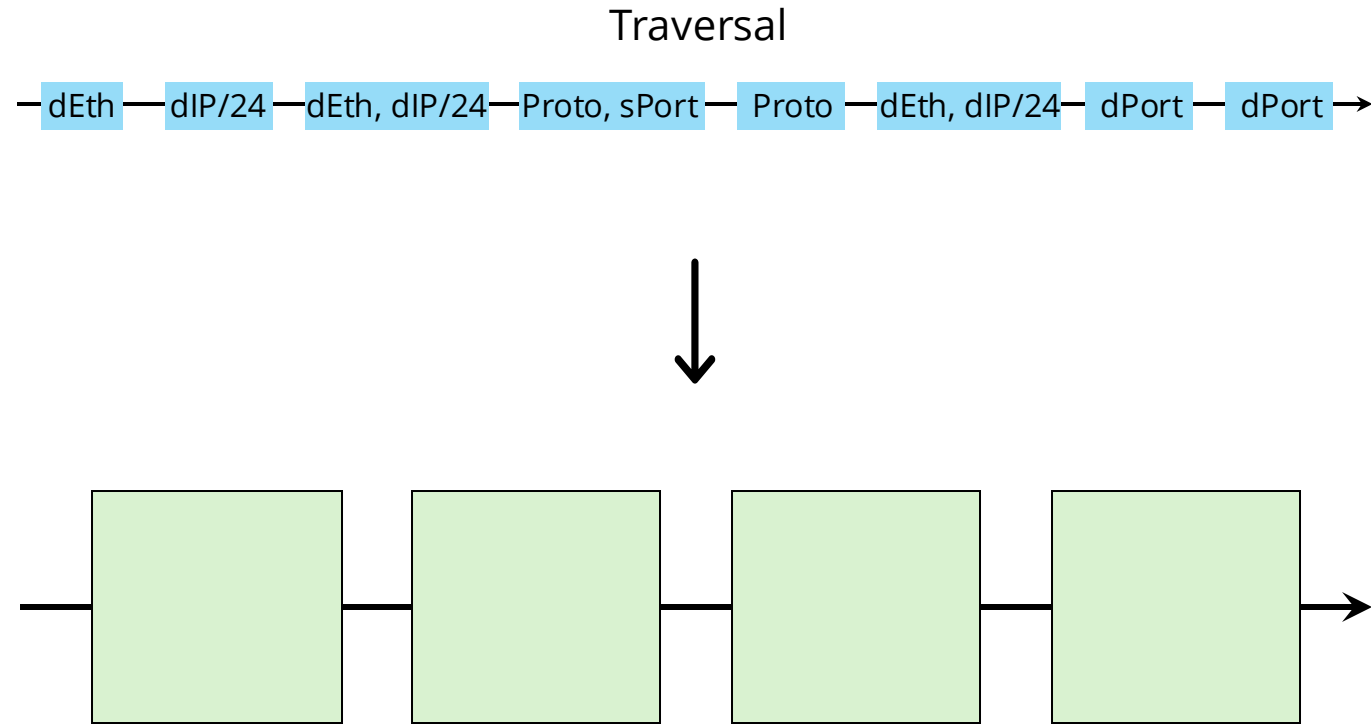
SmartNIC pipelines have a
handful of tables and offer
limited control flow



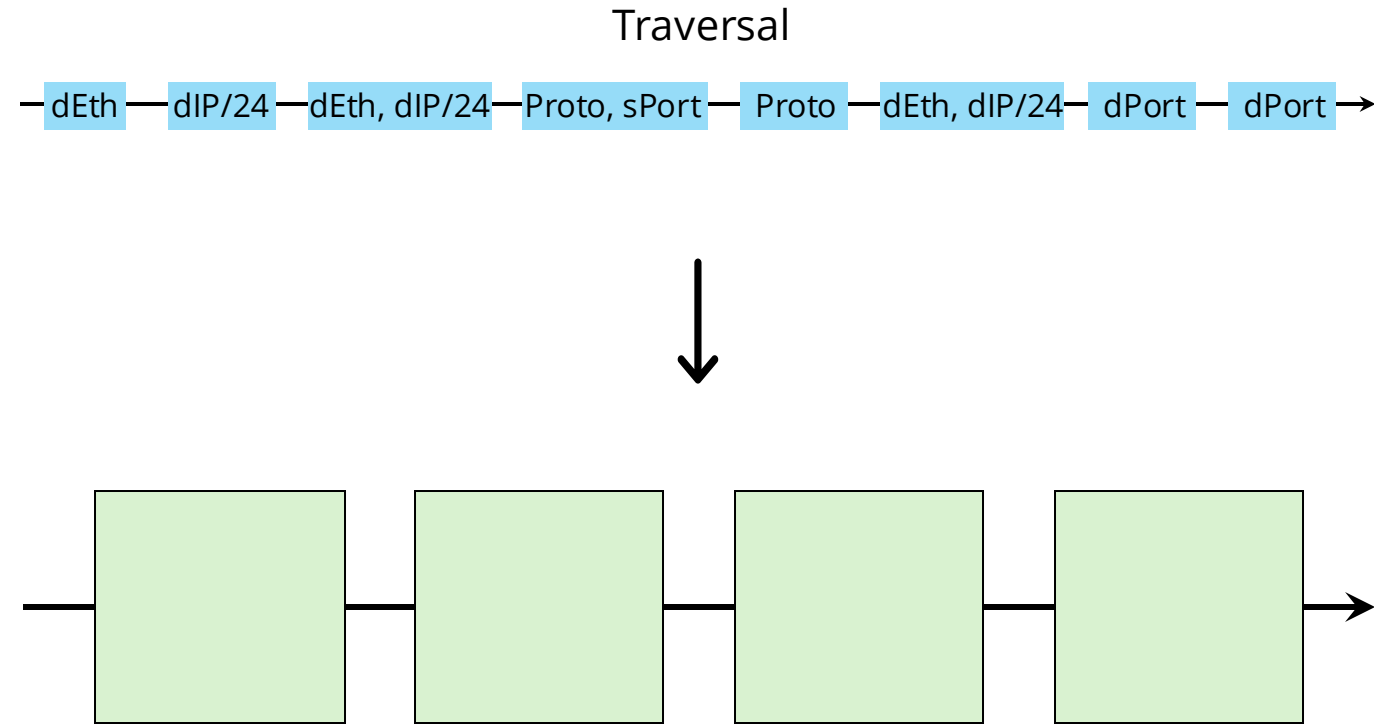
Cache Localities with Multiple Tables!



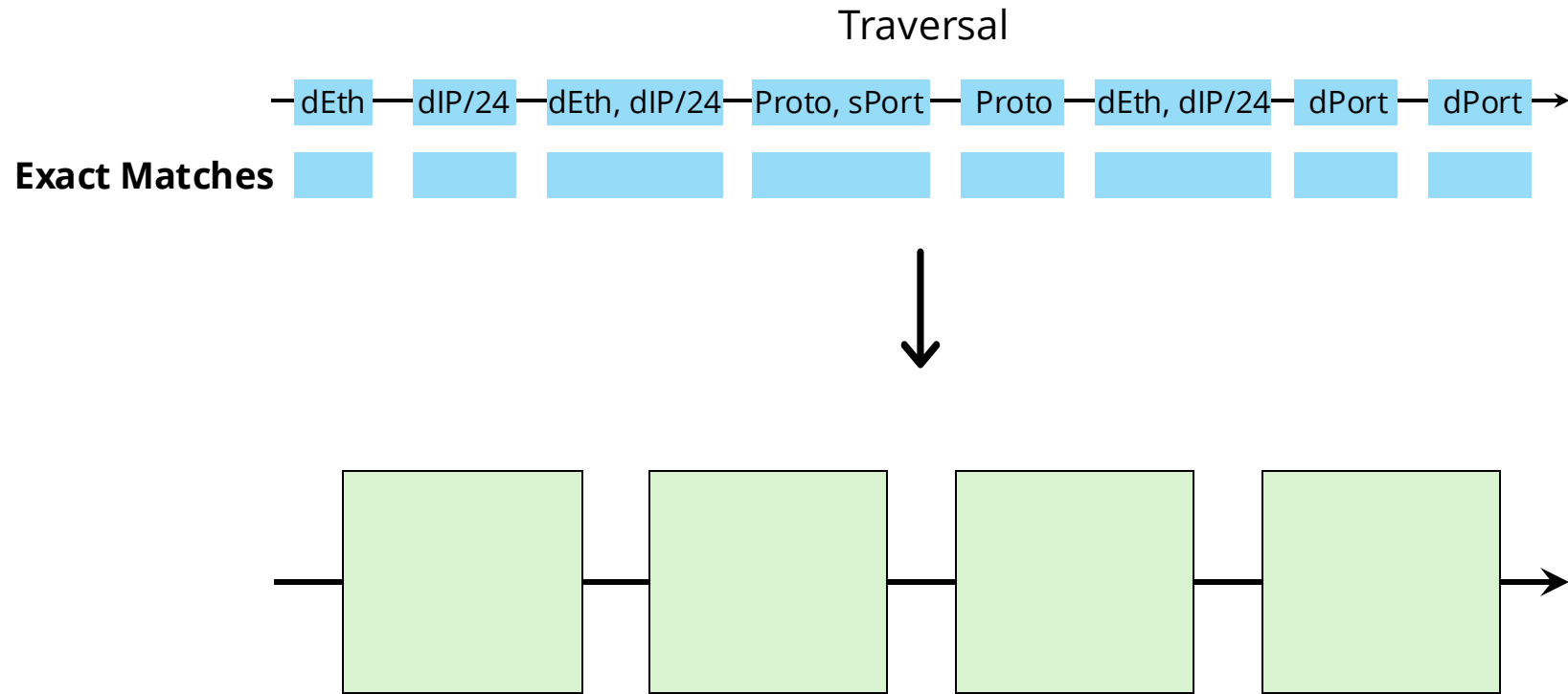
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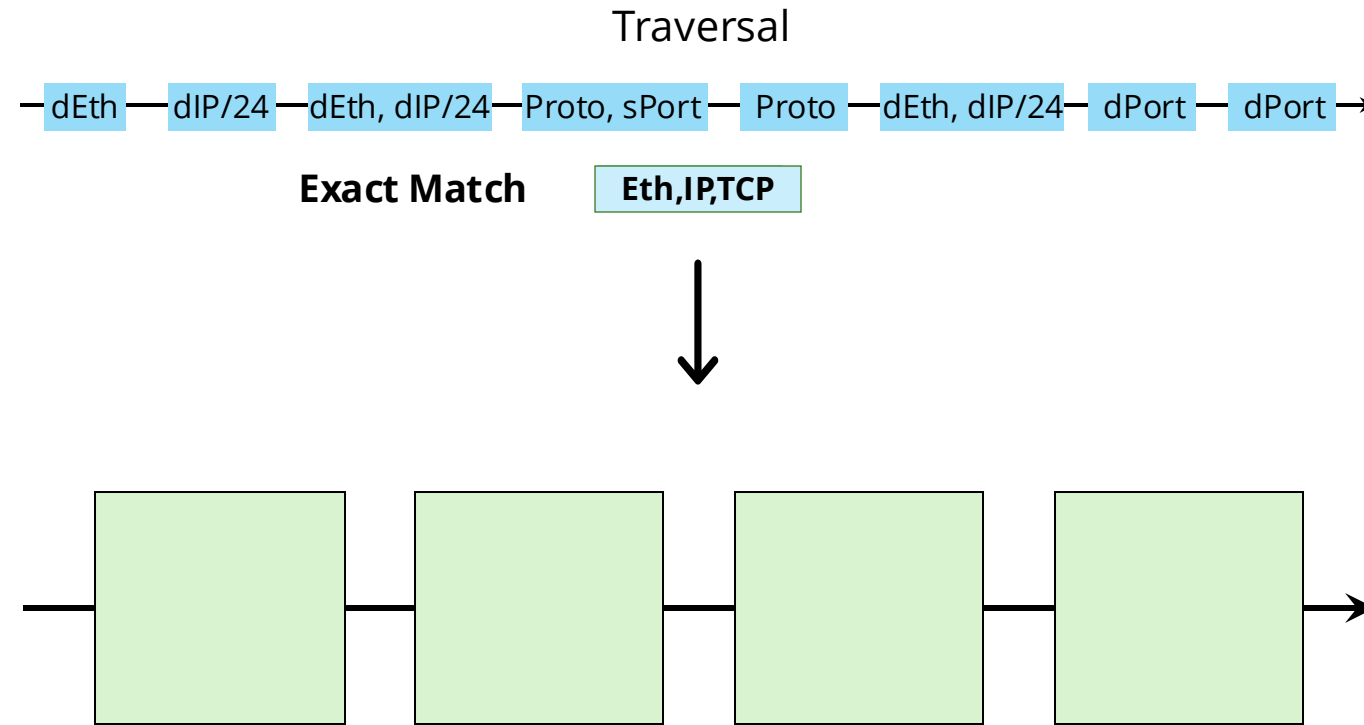
Caching *Temporal Locality* in the SmartNIC



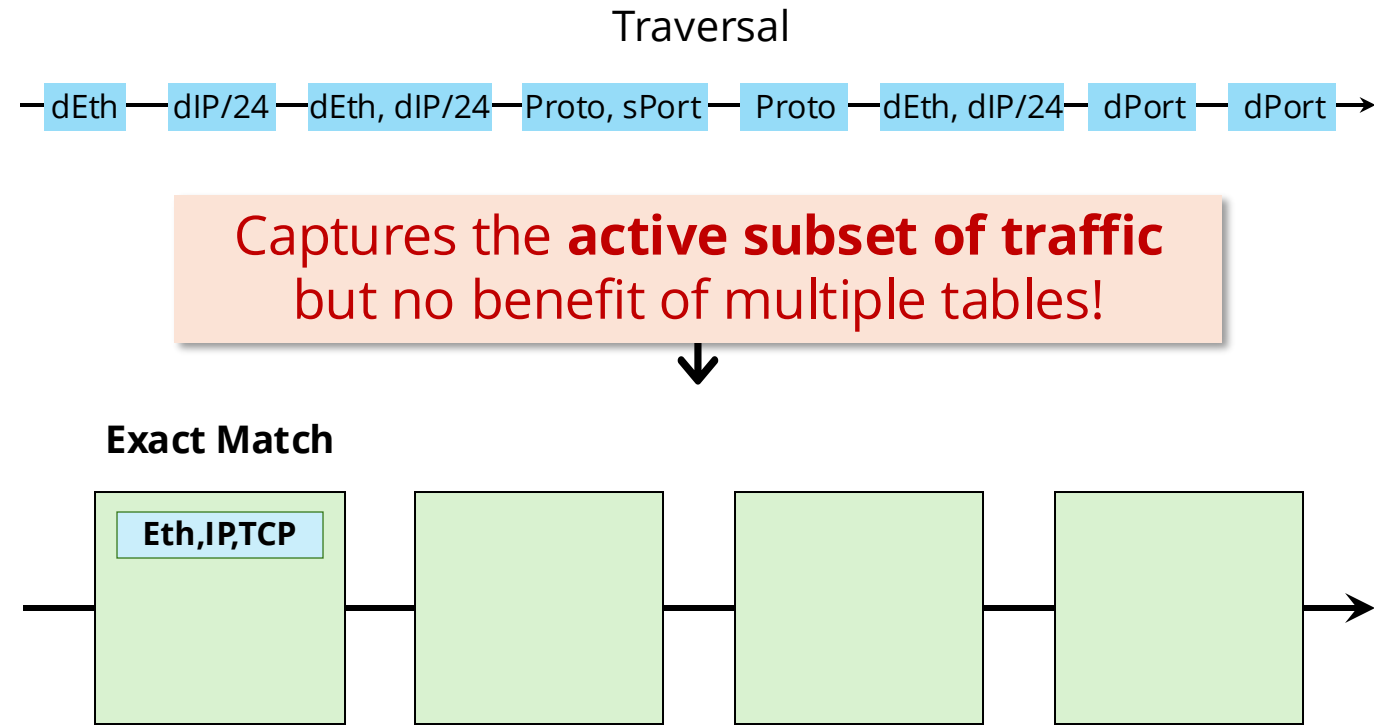
Caching *Temporal Locality* in the SmartNIC



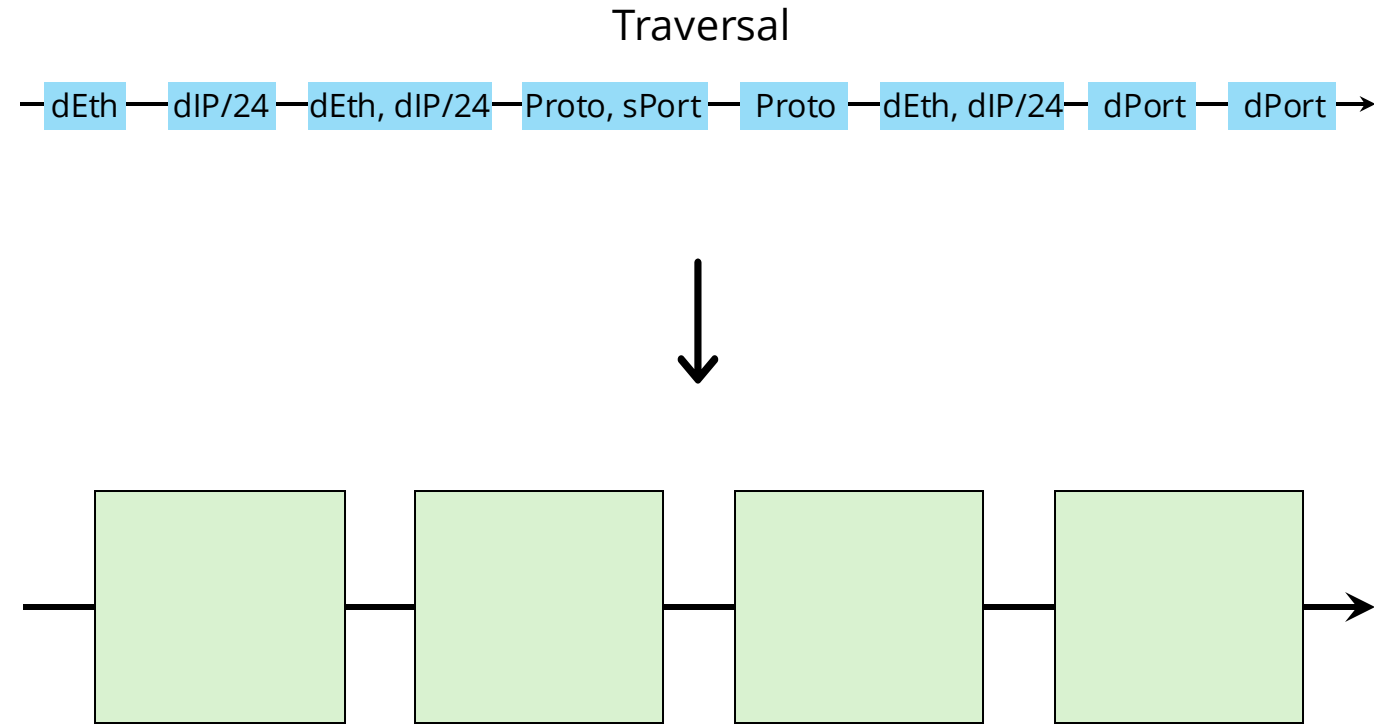
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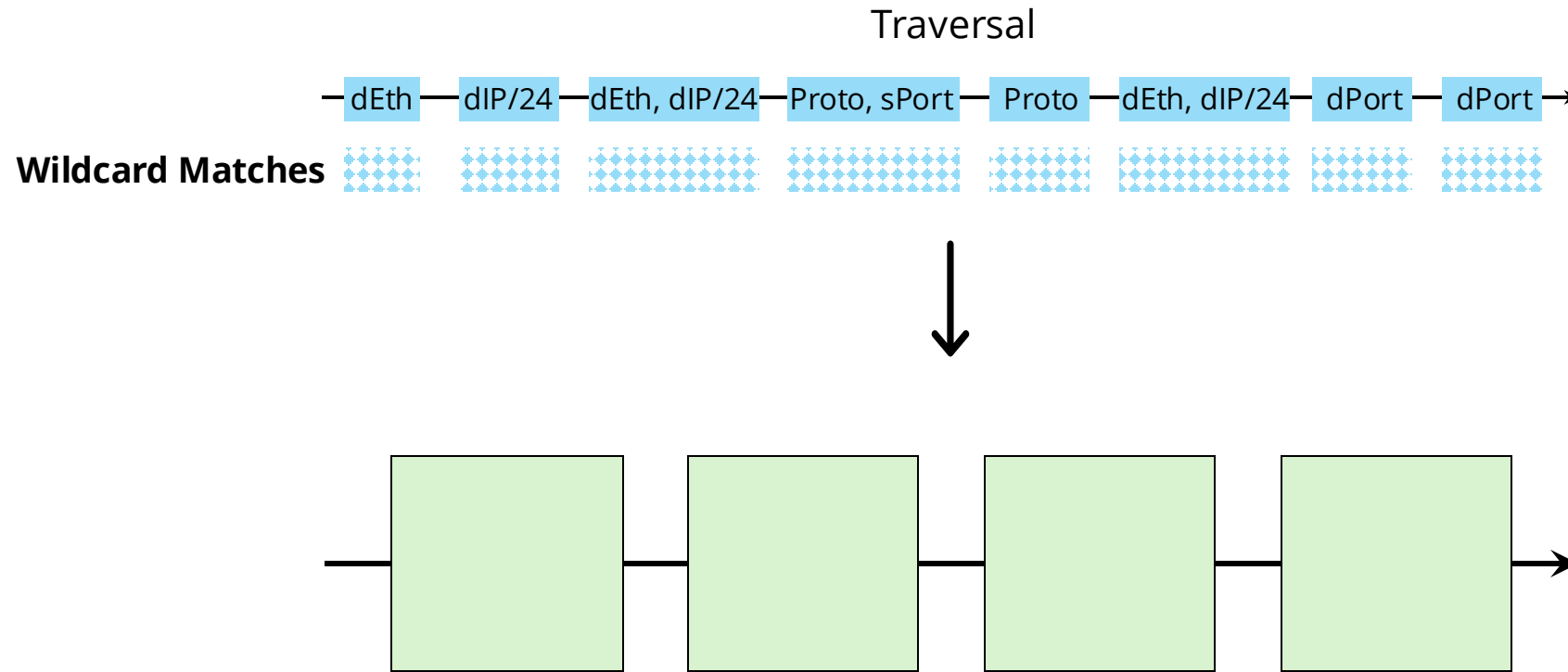
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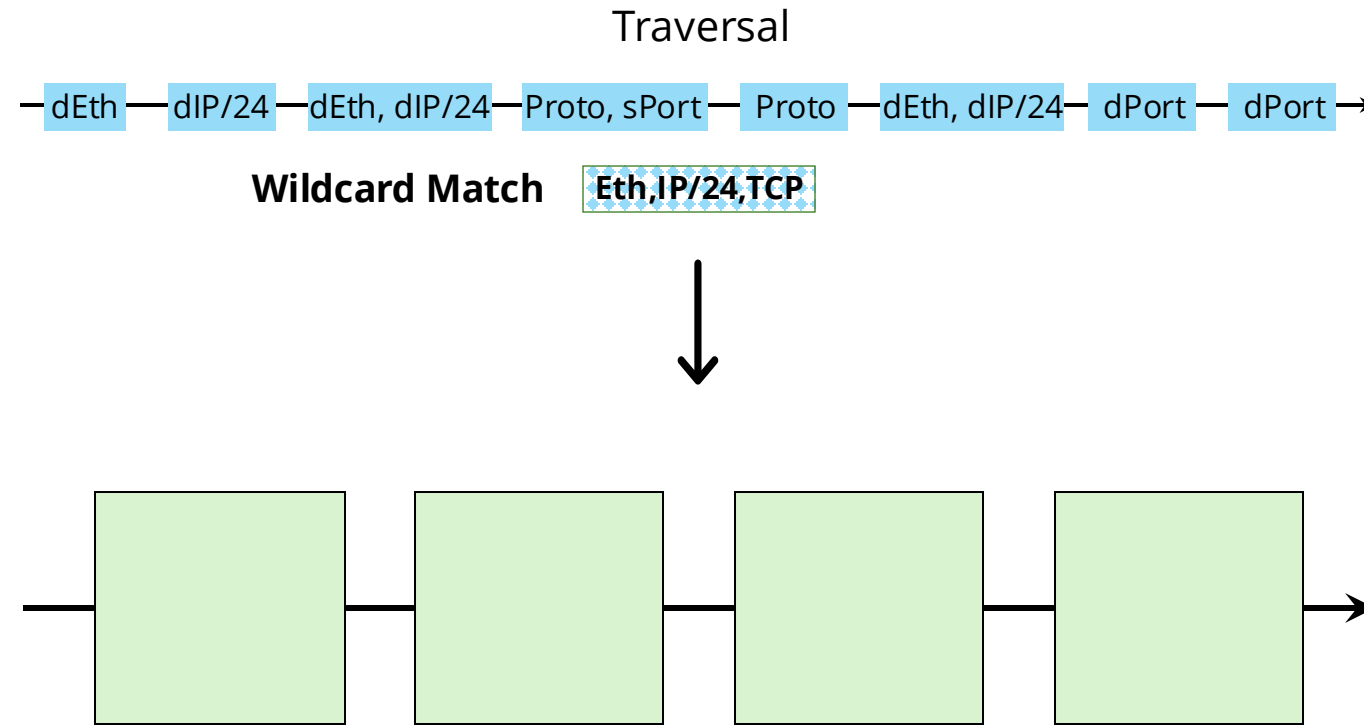
Caching *Spatial Locality* in the SmartNIC



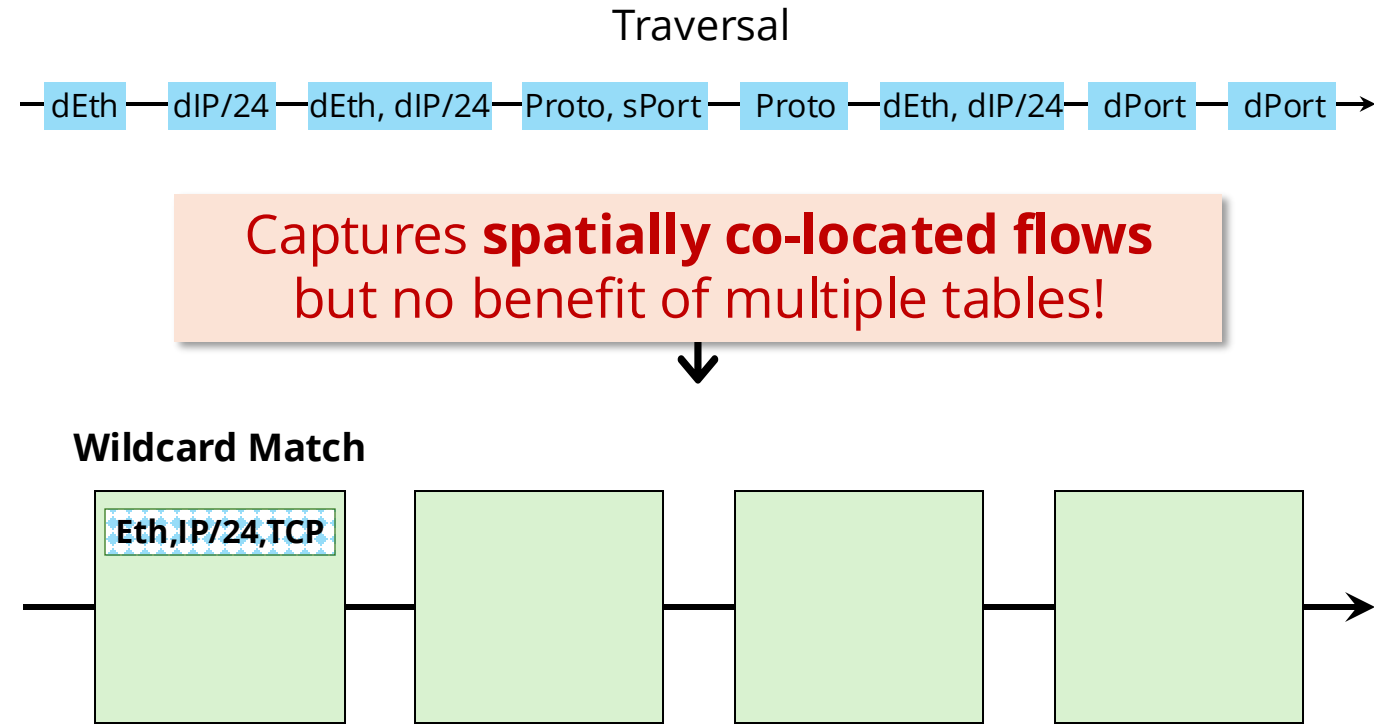
Caching *Spatial Locality* in the SmartNIC



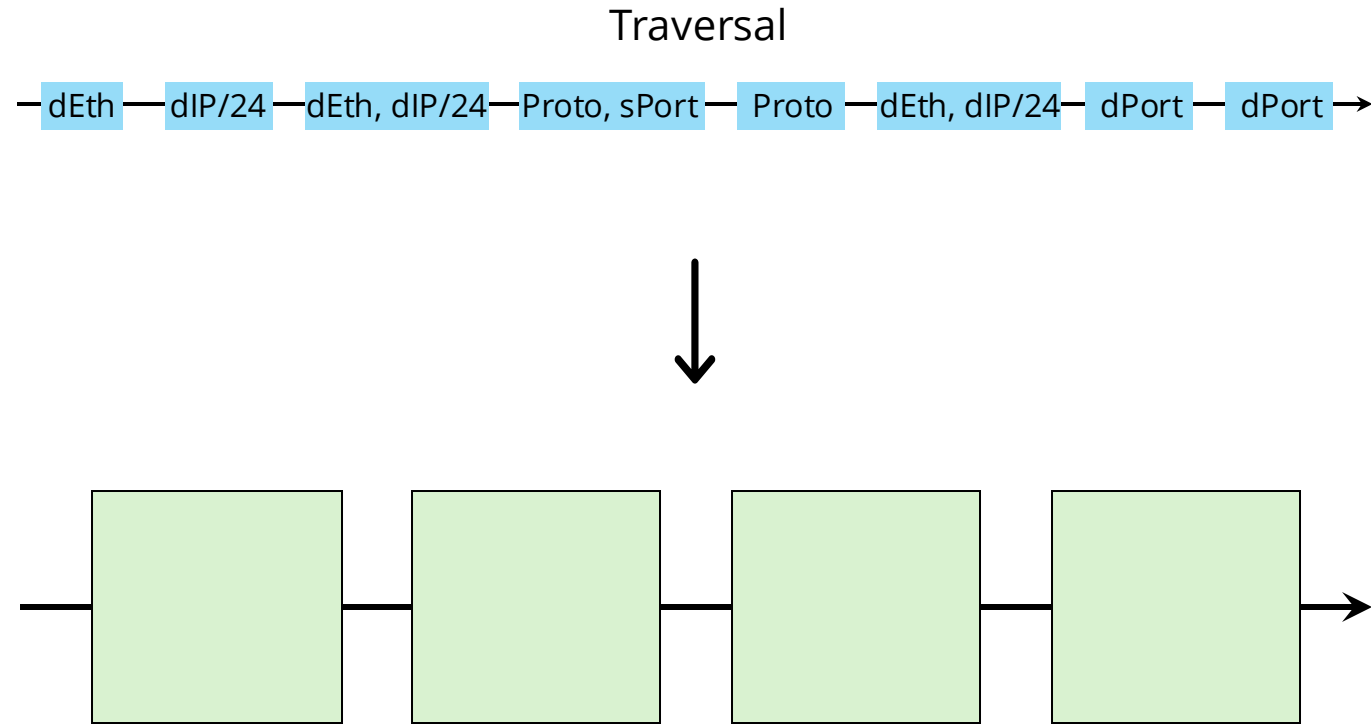
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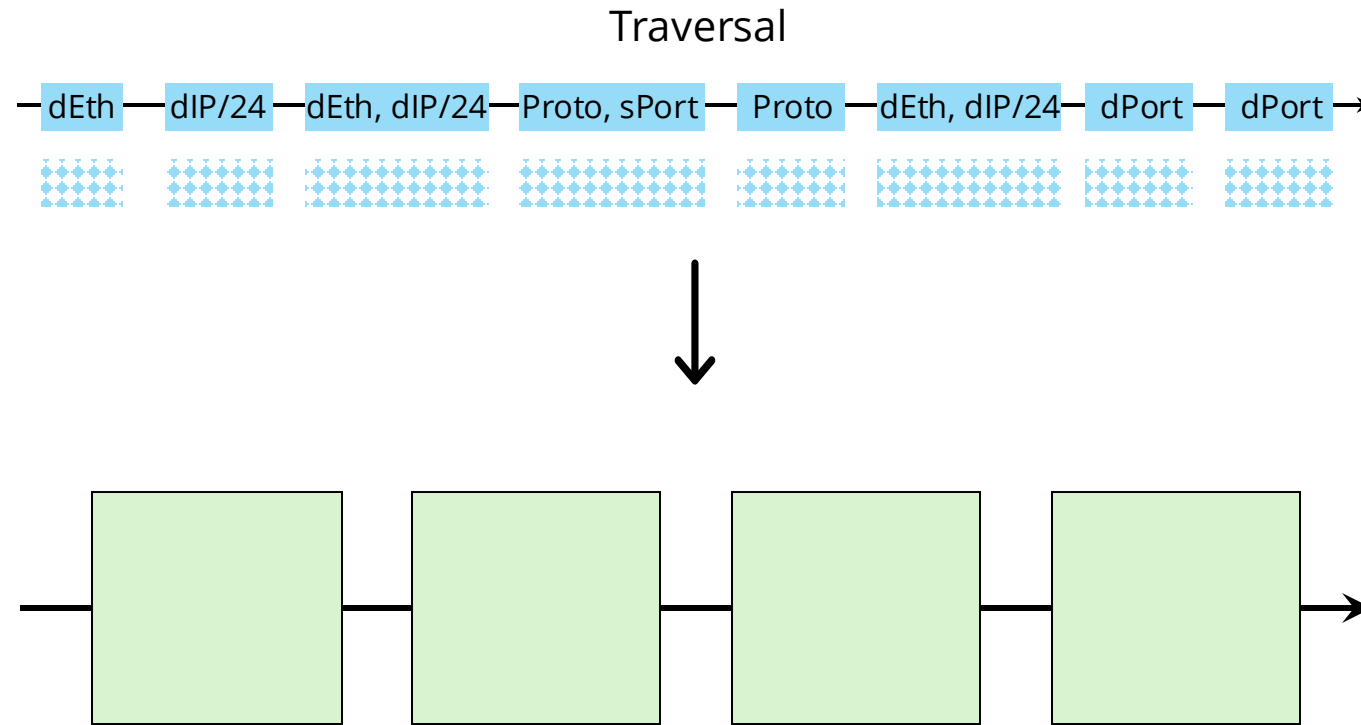
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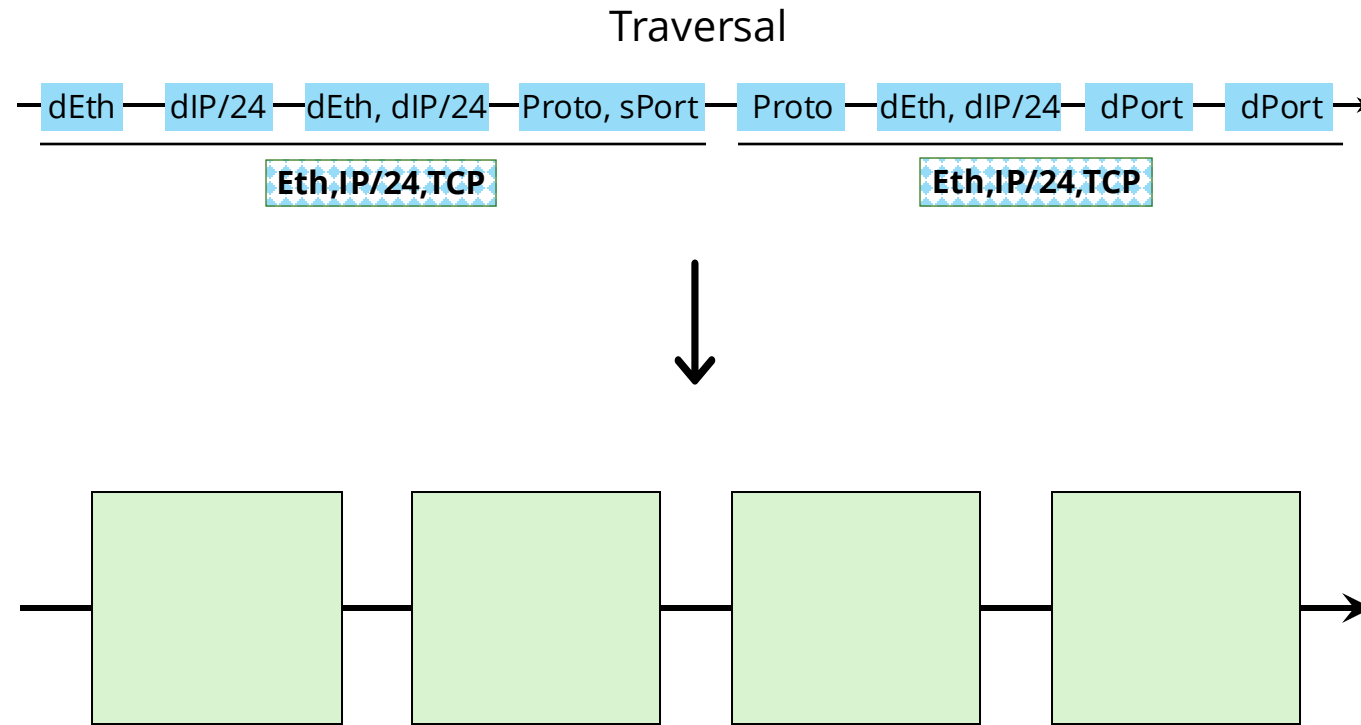
How about *Partitioning* the Traversal?



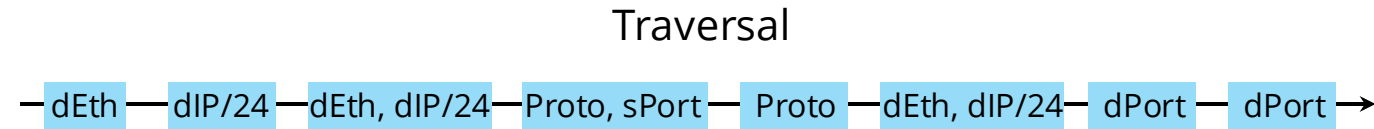
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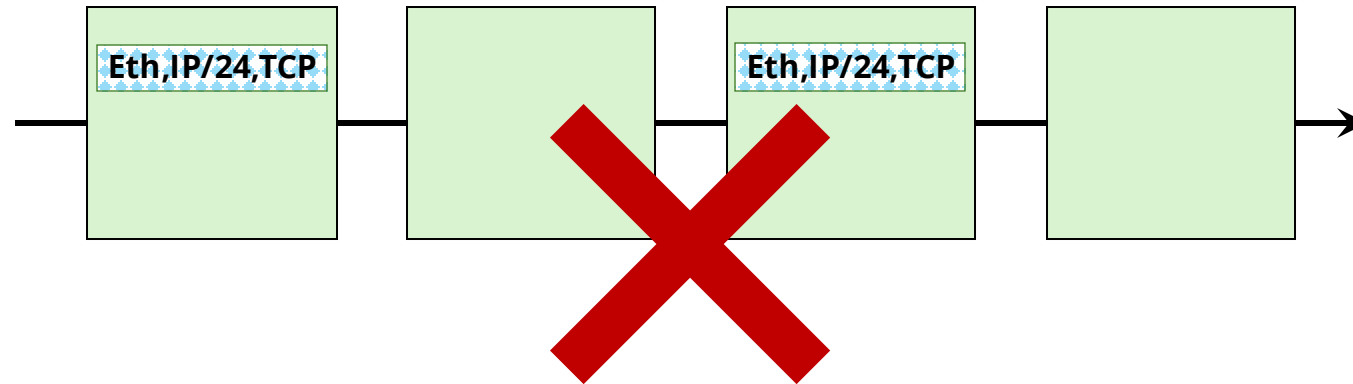
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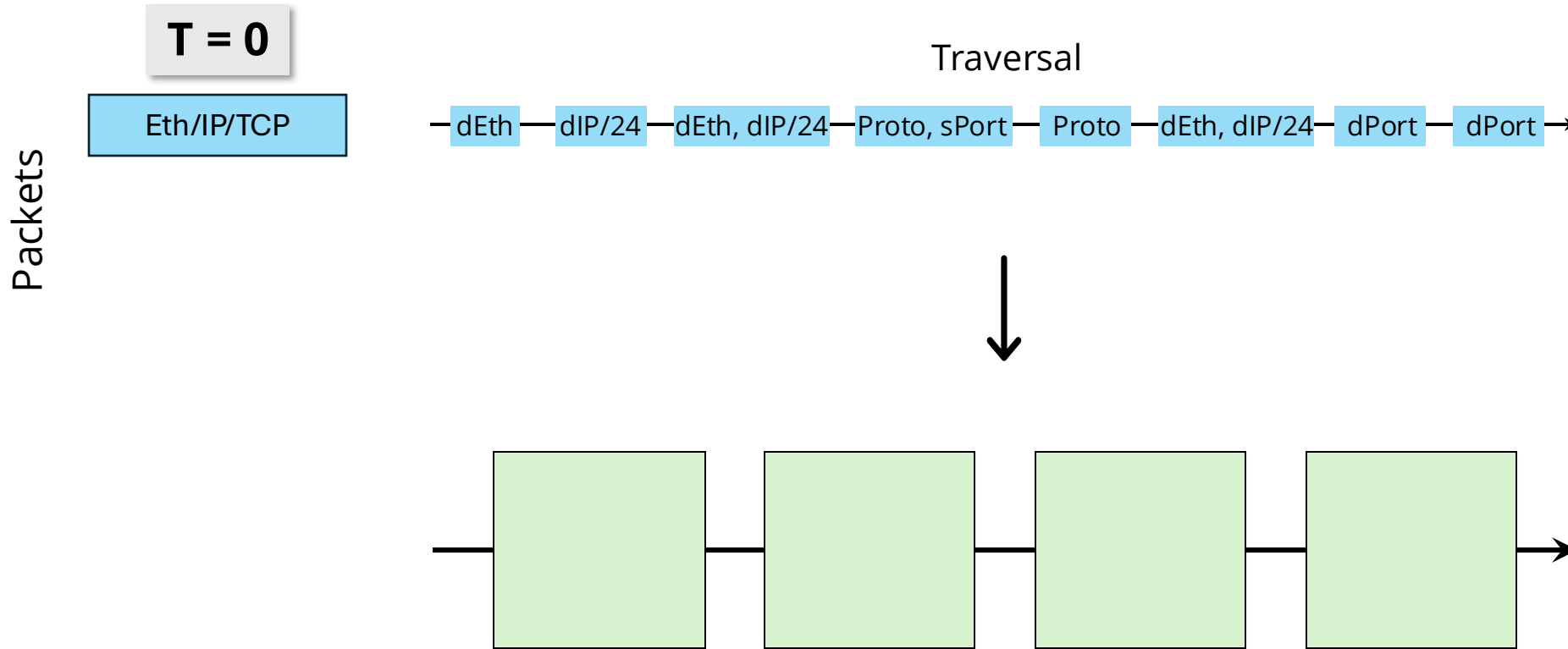
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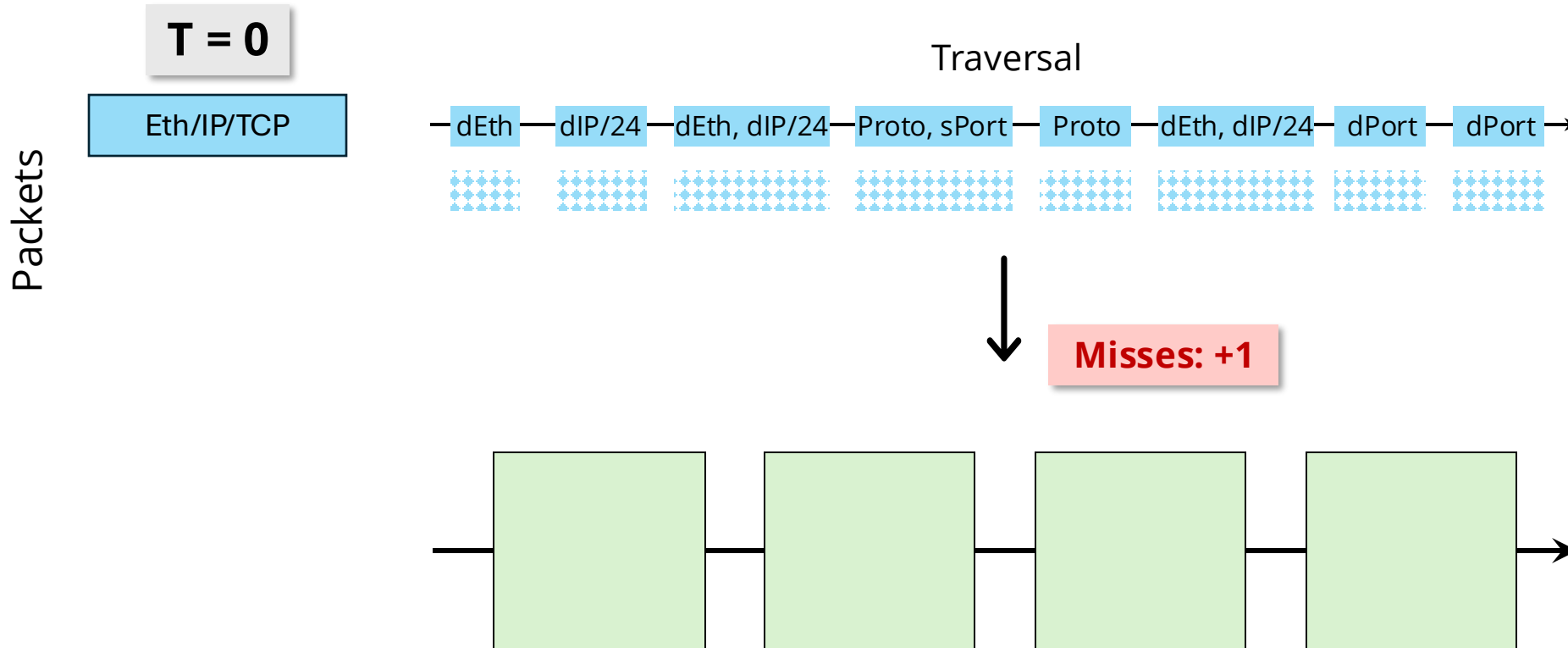
Arbitrary partitions can generate
identical cache entries
that **exhaust the cache** space faster!



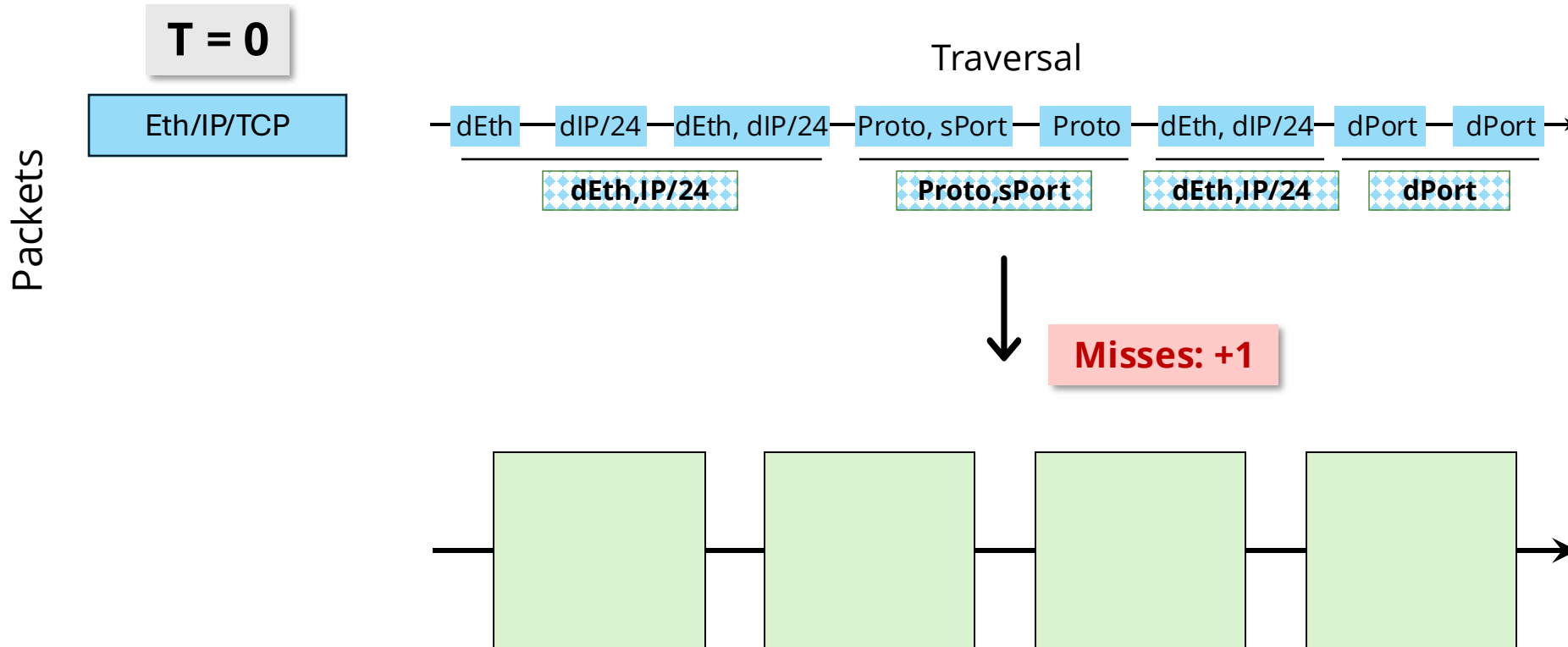
How about Smartly *Partitioning* the Traversal?



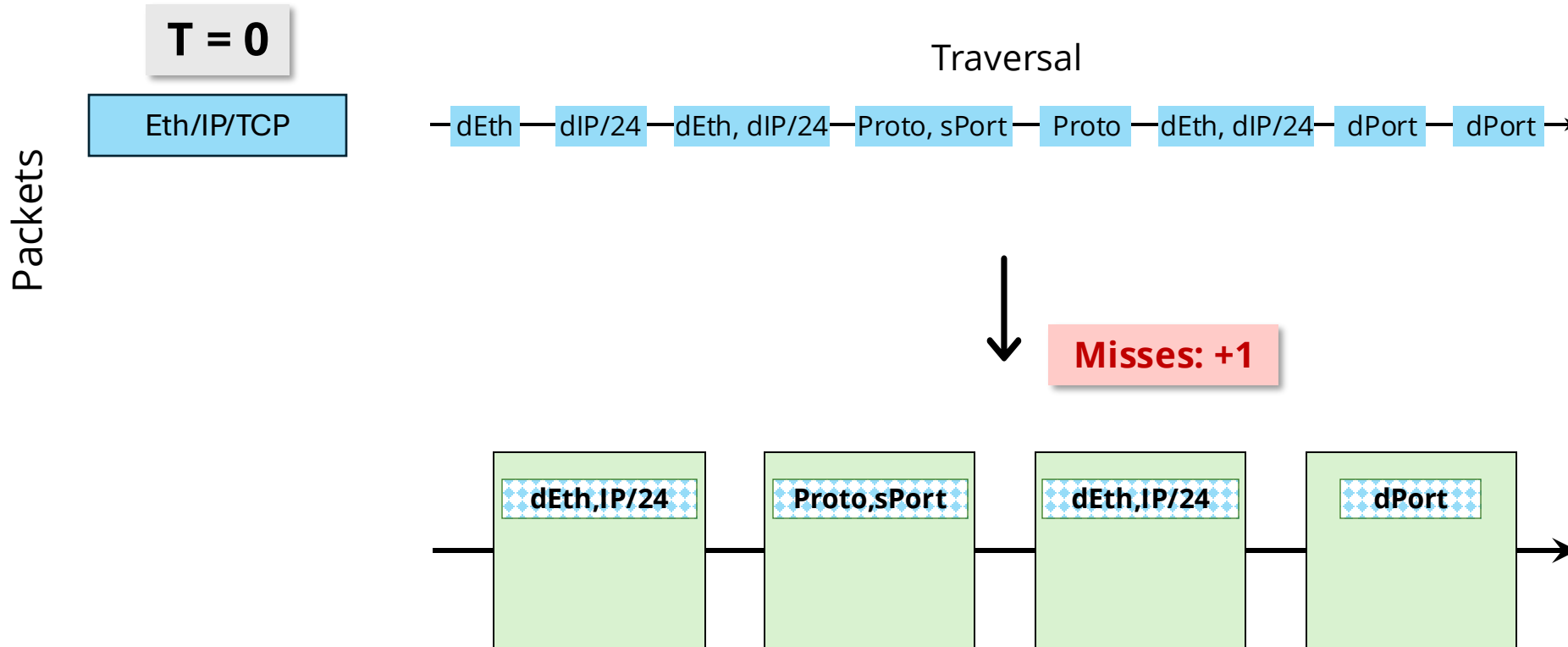
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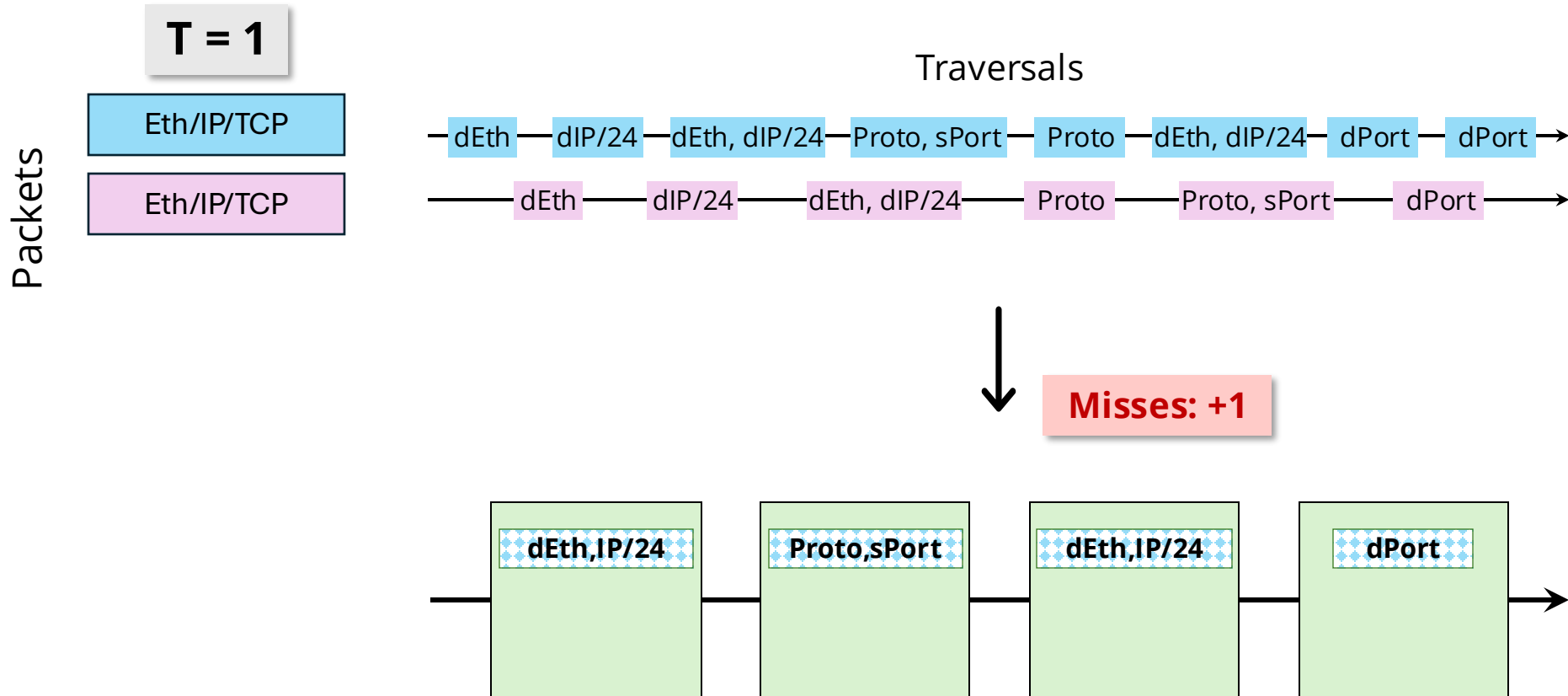
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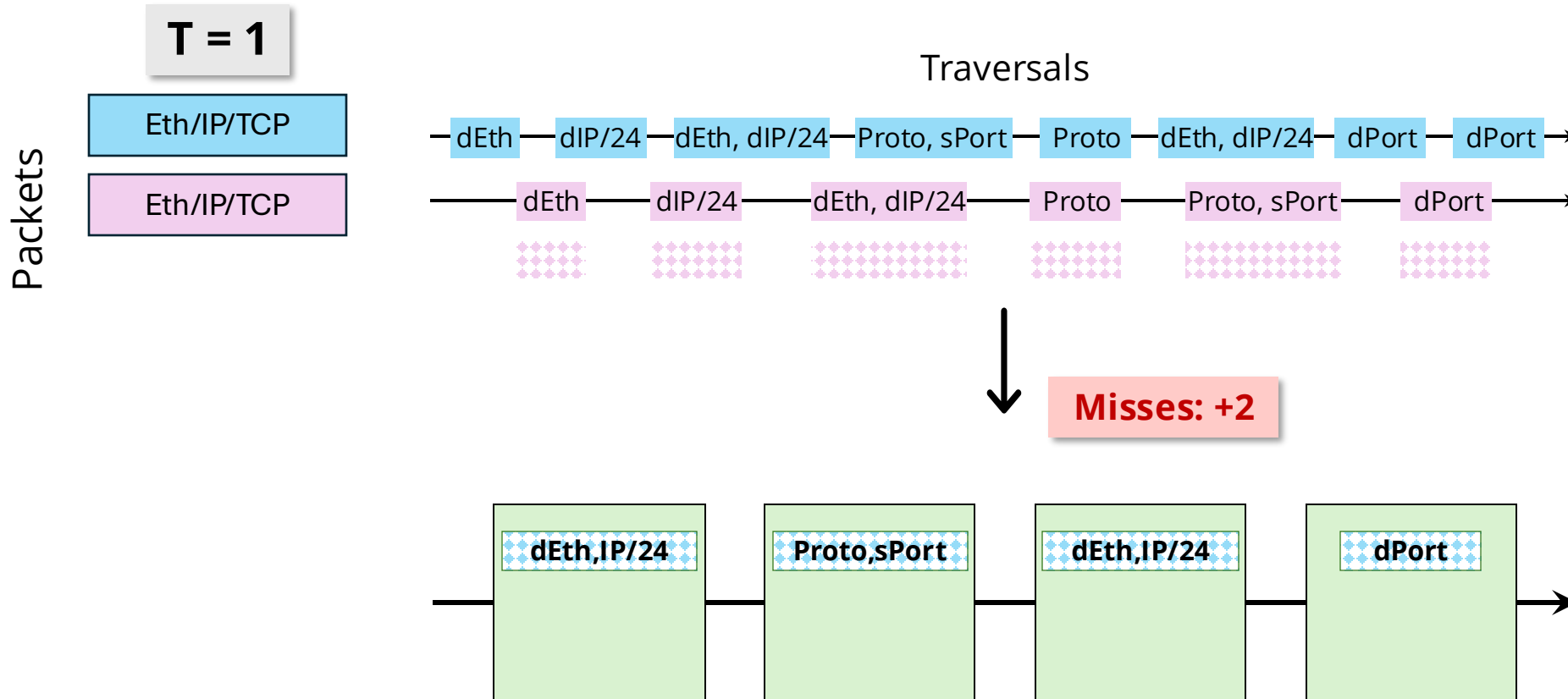
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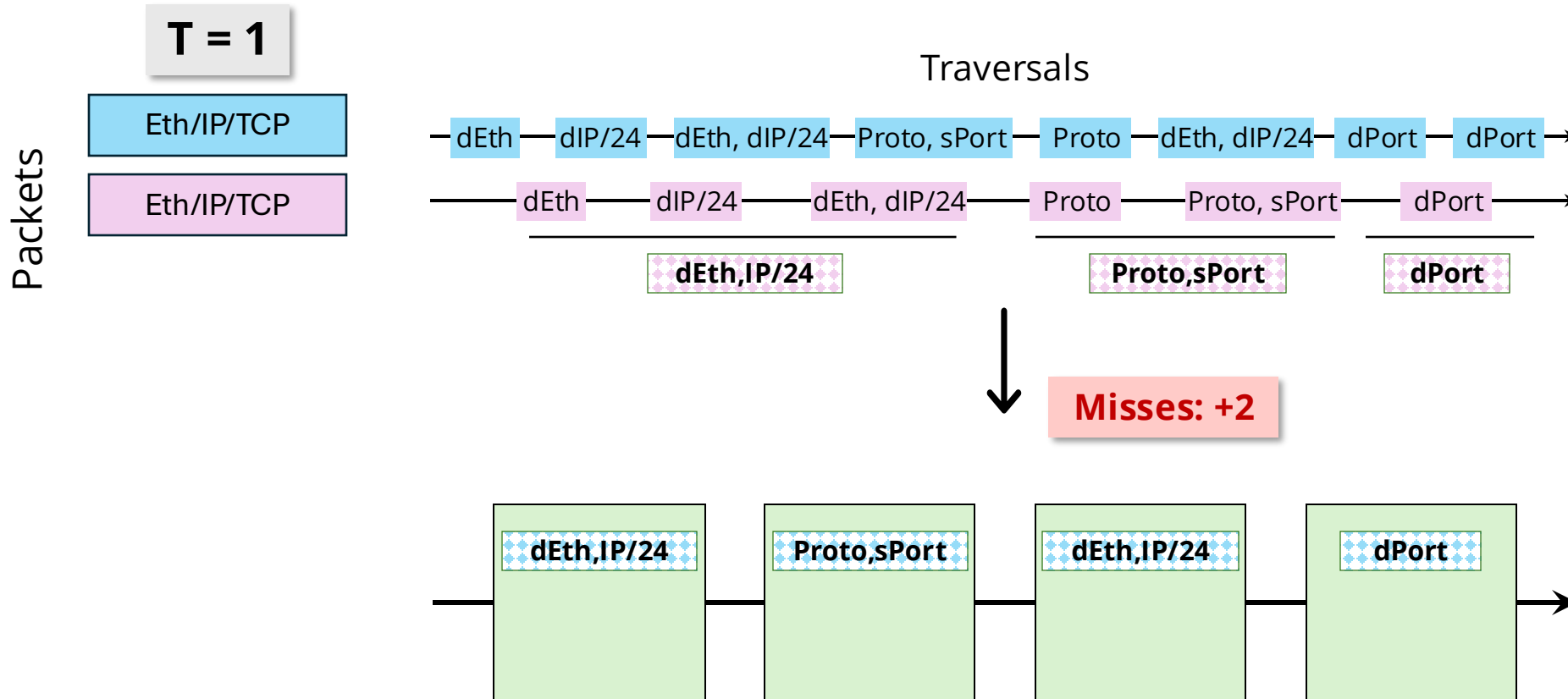
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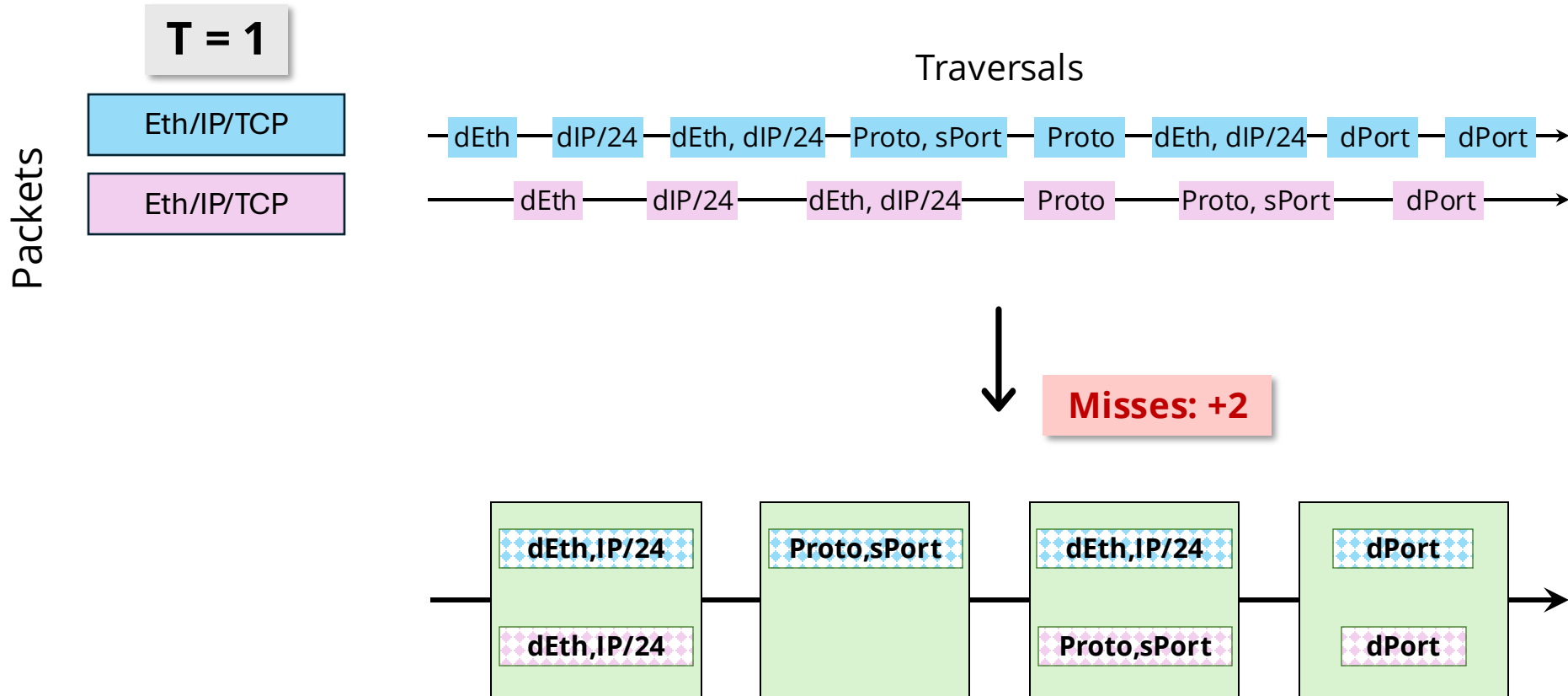
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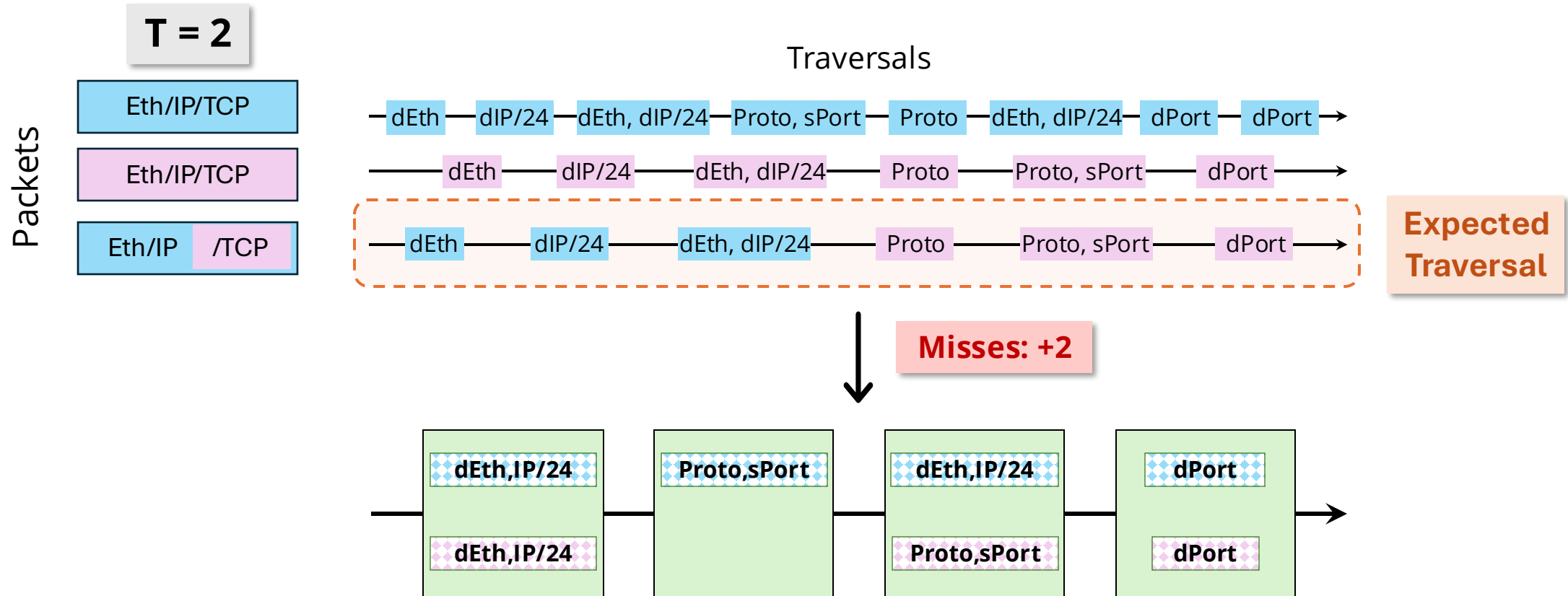
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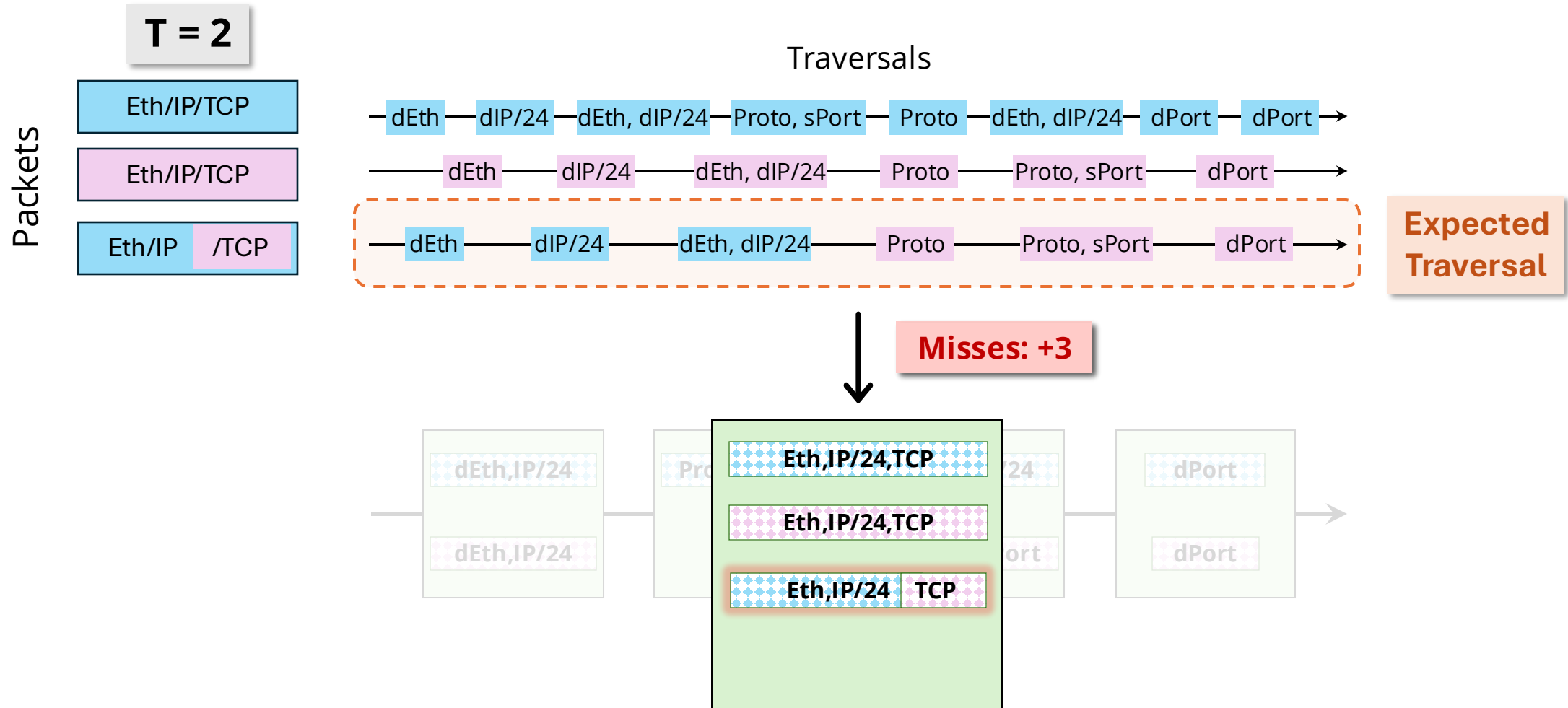
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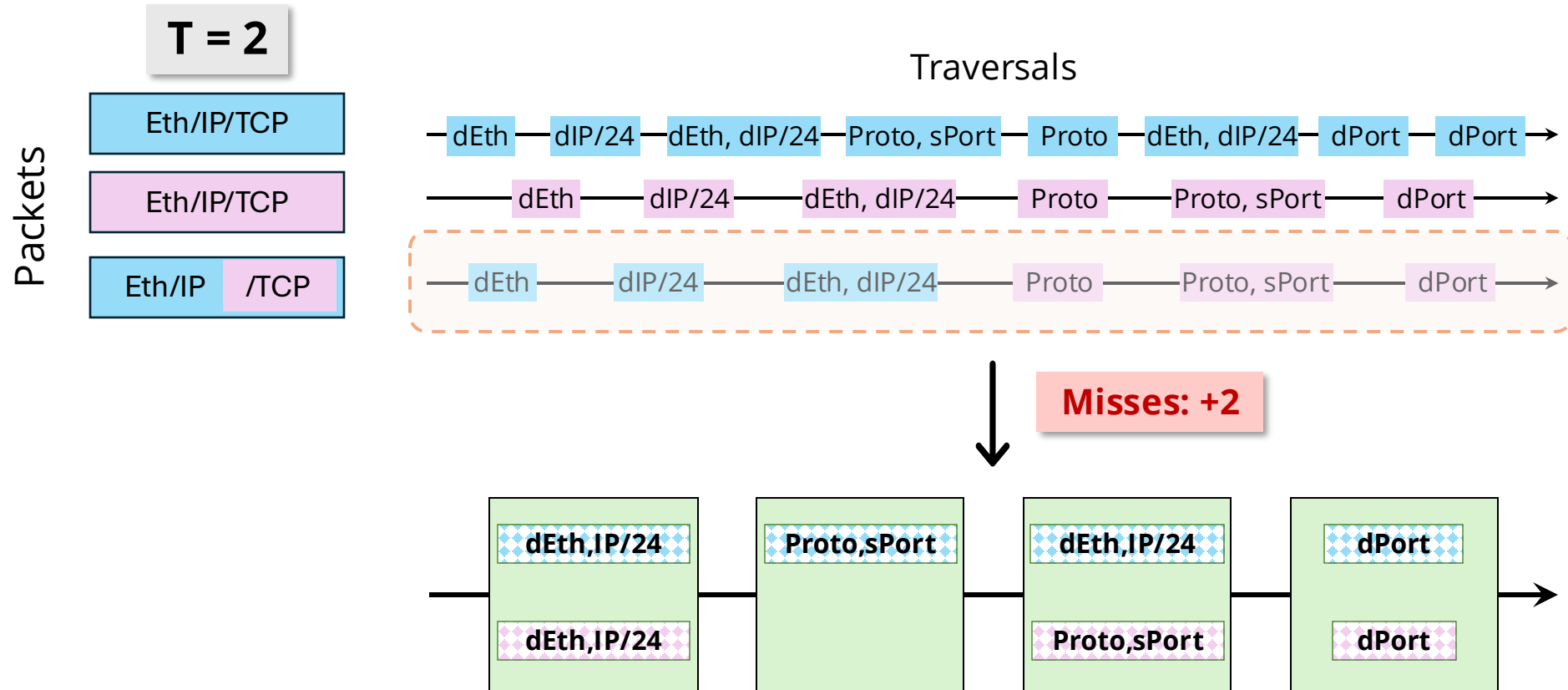
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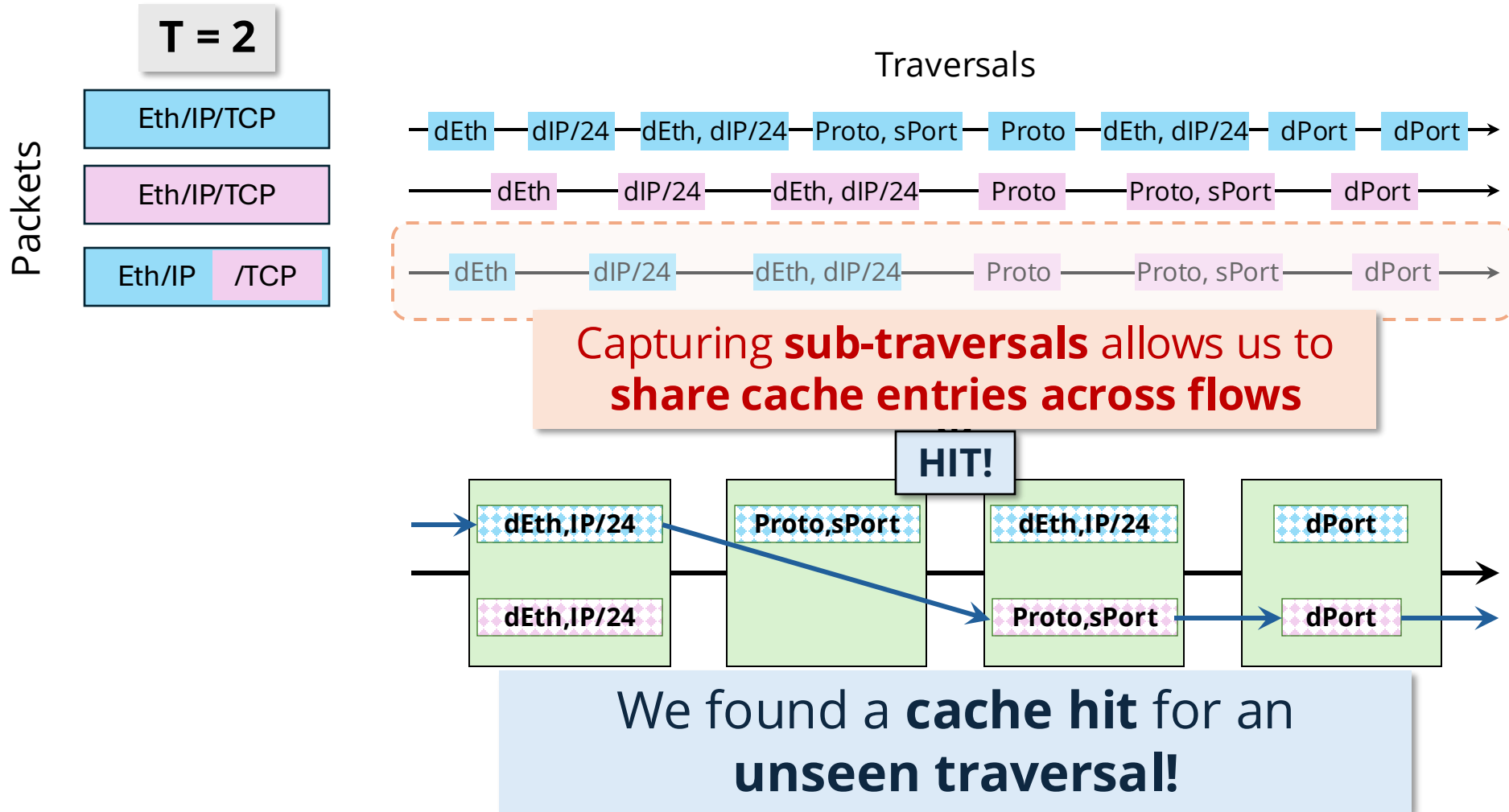
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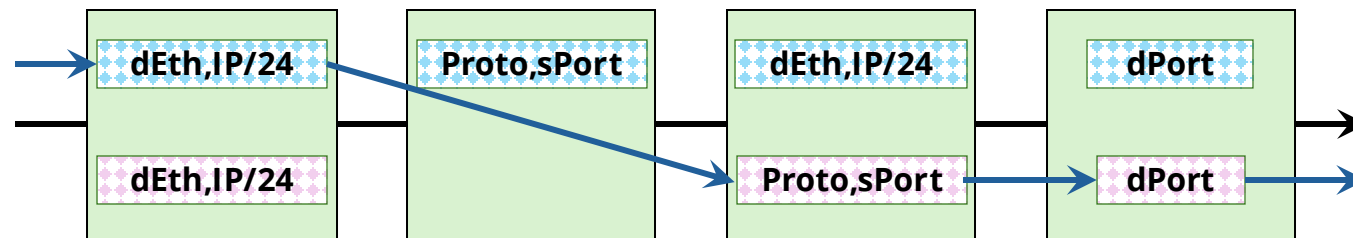
Sub-Traversal Caching with Gigaflow

Sub-traversals shared across flows
as cache entries in the SmartNIC

Optimal sub-traversals as cache entries
capture **pipeline-aware locality**



Up to **two orders of magnitude**
improvement in rule space coverage!



Sub-Traversal Caching with Gigaflow

Optimal sub-traversals as cache entries capture **pipeline-aware locality**



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SmartNIC Cache	Real-World vSwitch Pipelines (Ps)				
	P1	P2	P3	P4	P5
MF	32K	32K	32K	32K	32K
GF	??				

Rule Space Coverage of Megaflow 32K (MF)
vs Gigaflow 4x8K (GF) Cache

Capturing Pipeline-Aware Locality

Algorithm 1 Strawman Approach

Input: Set of traversals $T = \{t_1, t_2, \dots, t_N\}$

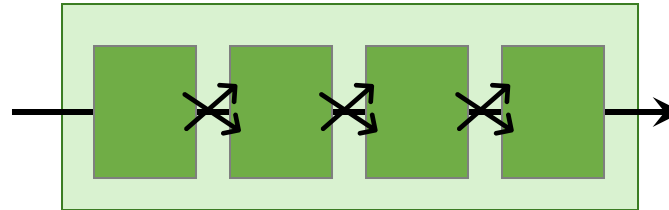
Output: Sub-traversals with maximum rule space coverage

This cost is **proportional** to **active flows**!
Keeps **increasing** with each **cache miss**!
Infeasible at line rate!

How to capture **pipeline-aware** locality
efficiently?

Capturing Pipeline-Aware Locality *Efficiently*

To goal is to **maximize** the **rule space coverage** in the SmartNIC tables



This **cross-product** rule space **should be maximized!**

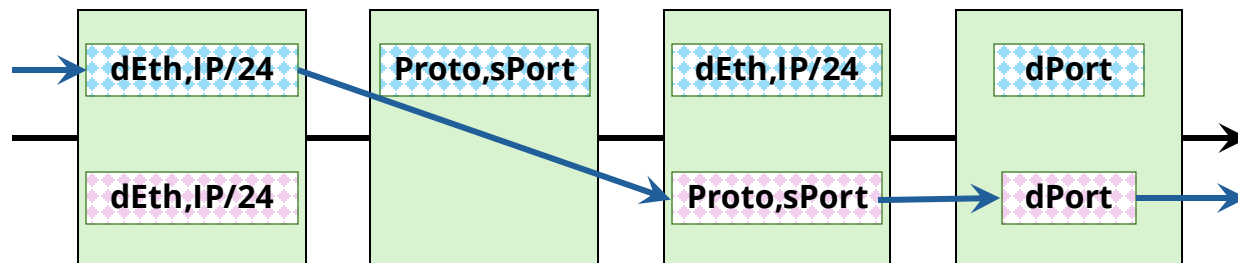


Pipeline-Aware locality → **Disjointedness**

Maximize the *field-level disjointedness*
among consecutive sub-traversals

Capturing Pipeline-Aware Locality *Efficiently*

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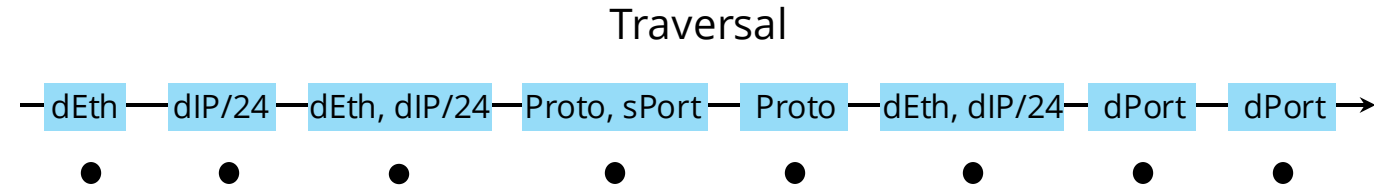
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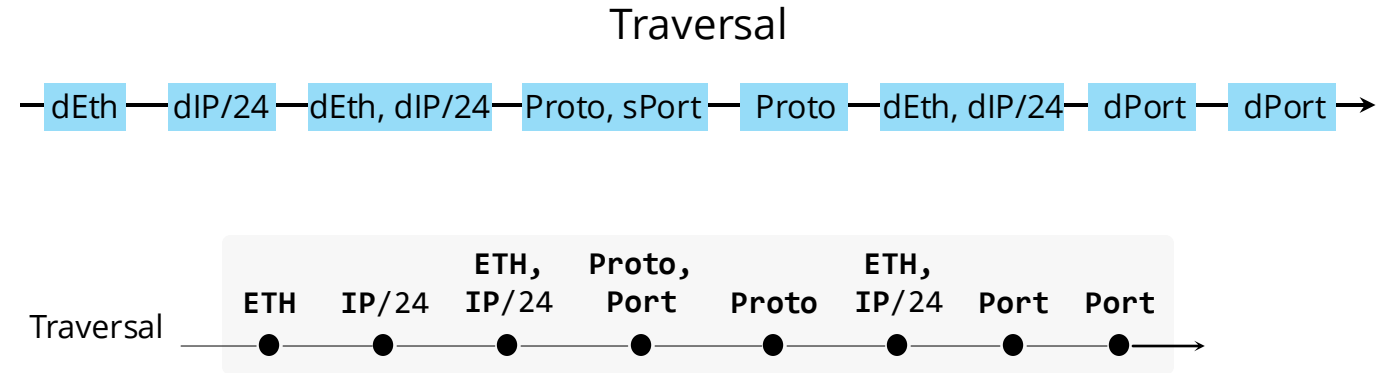
Disjointedness for Pipeline-Aware Locality



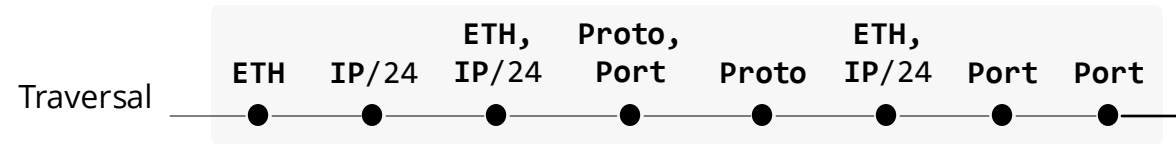
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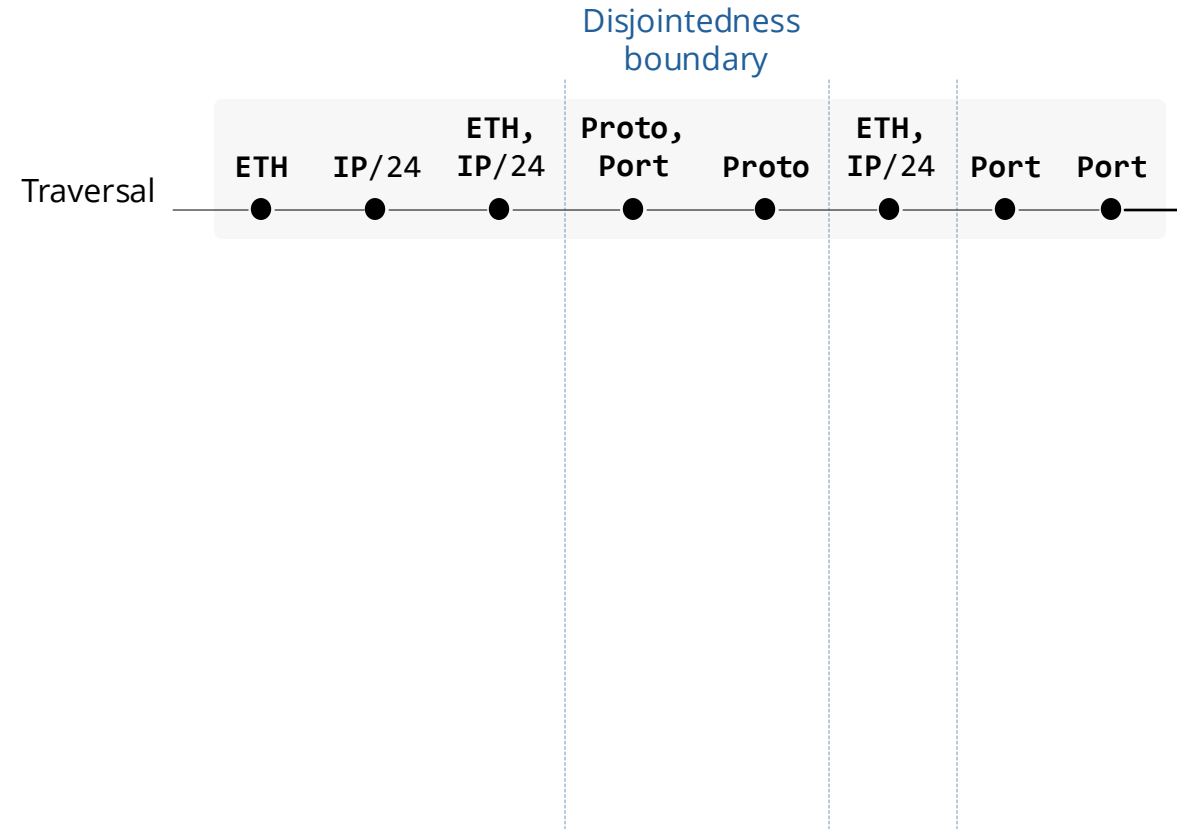
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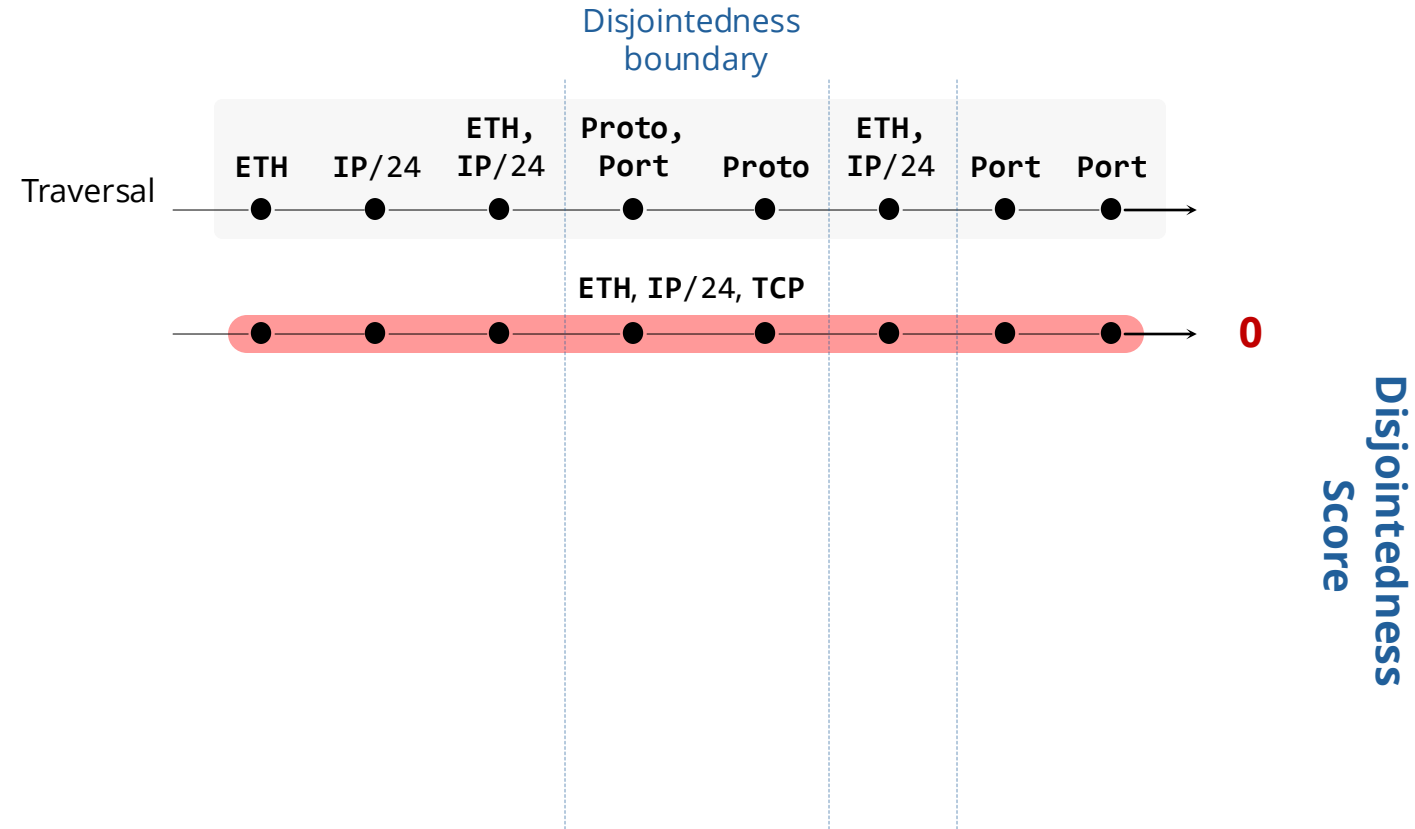
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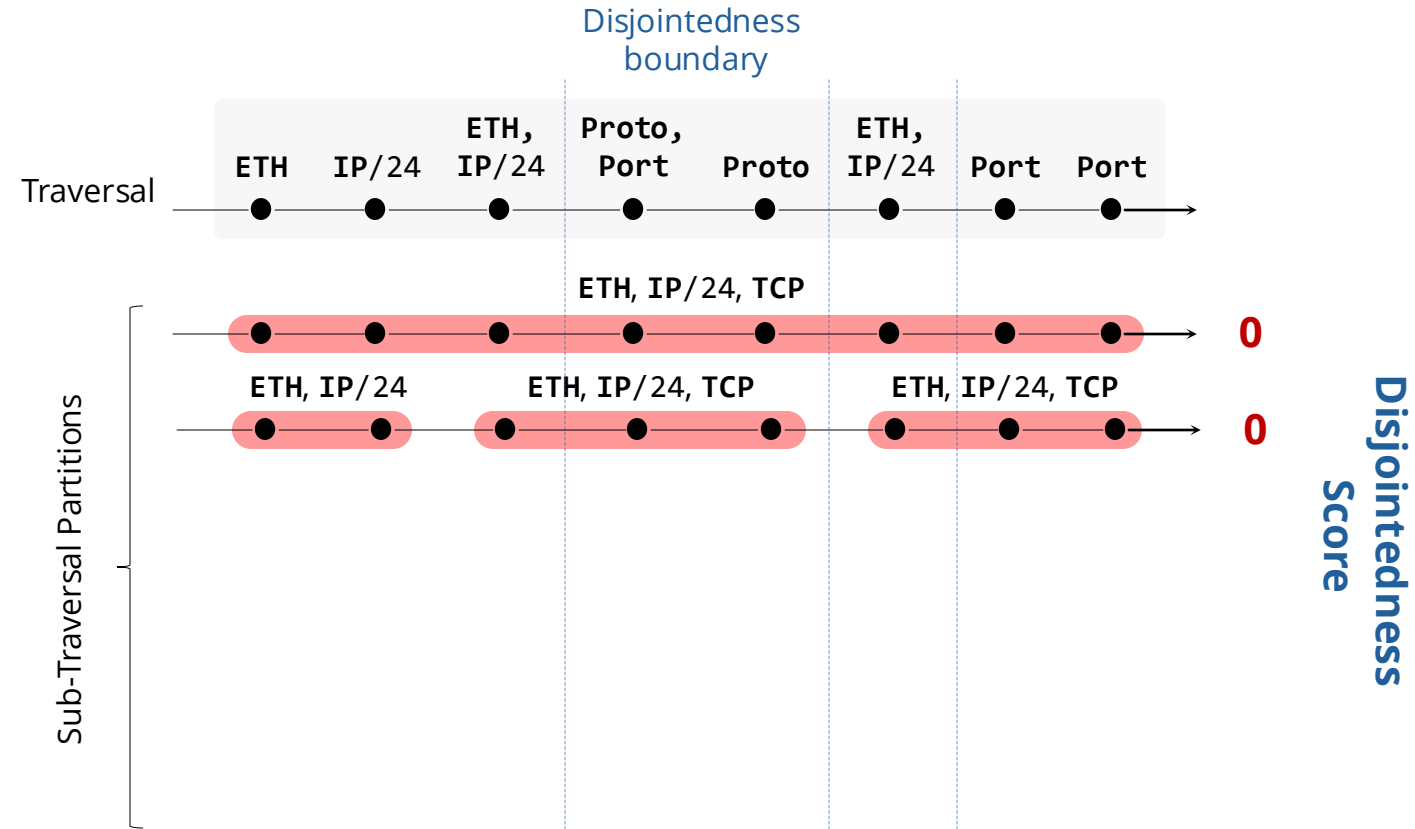
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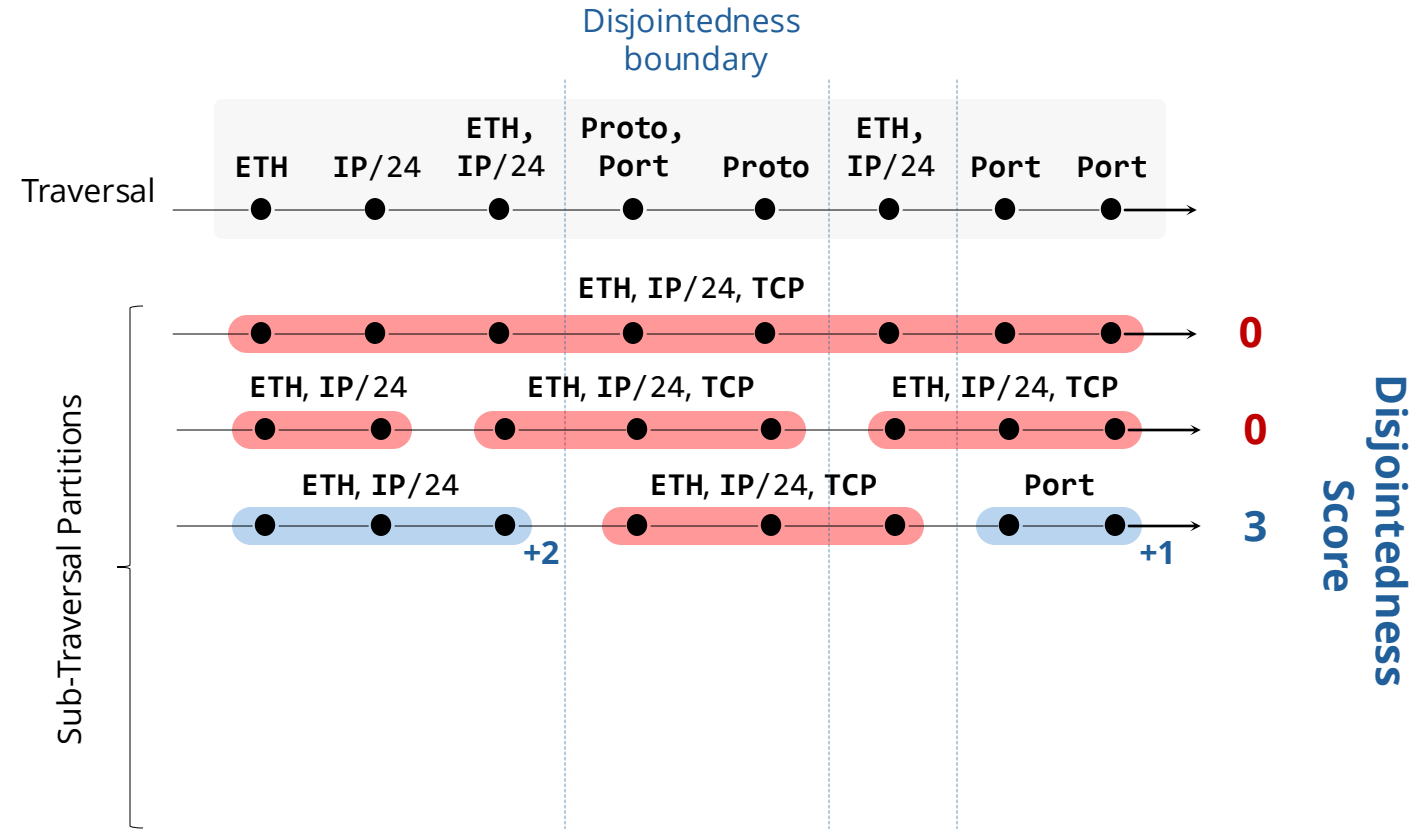
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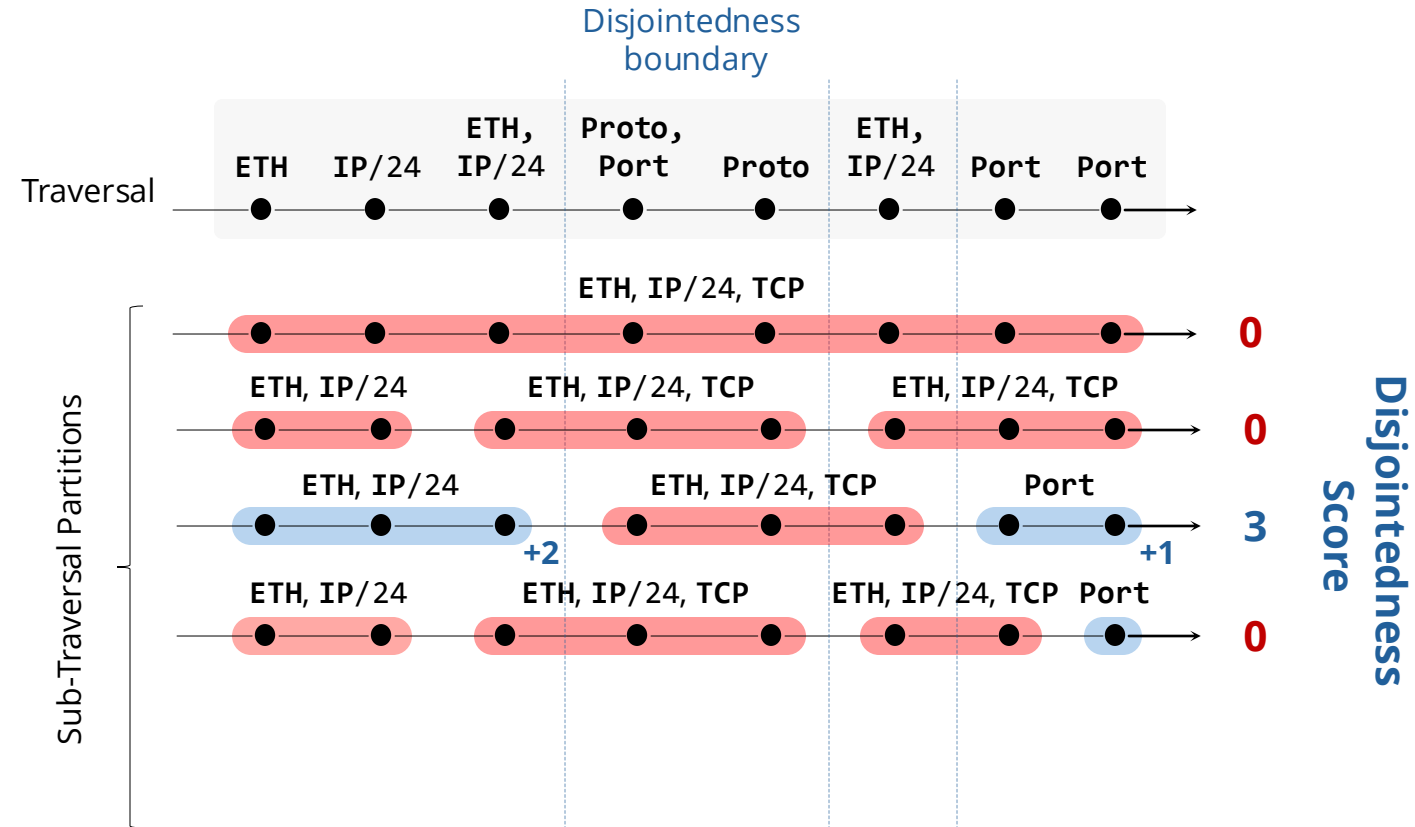
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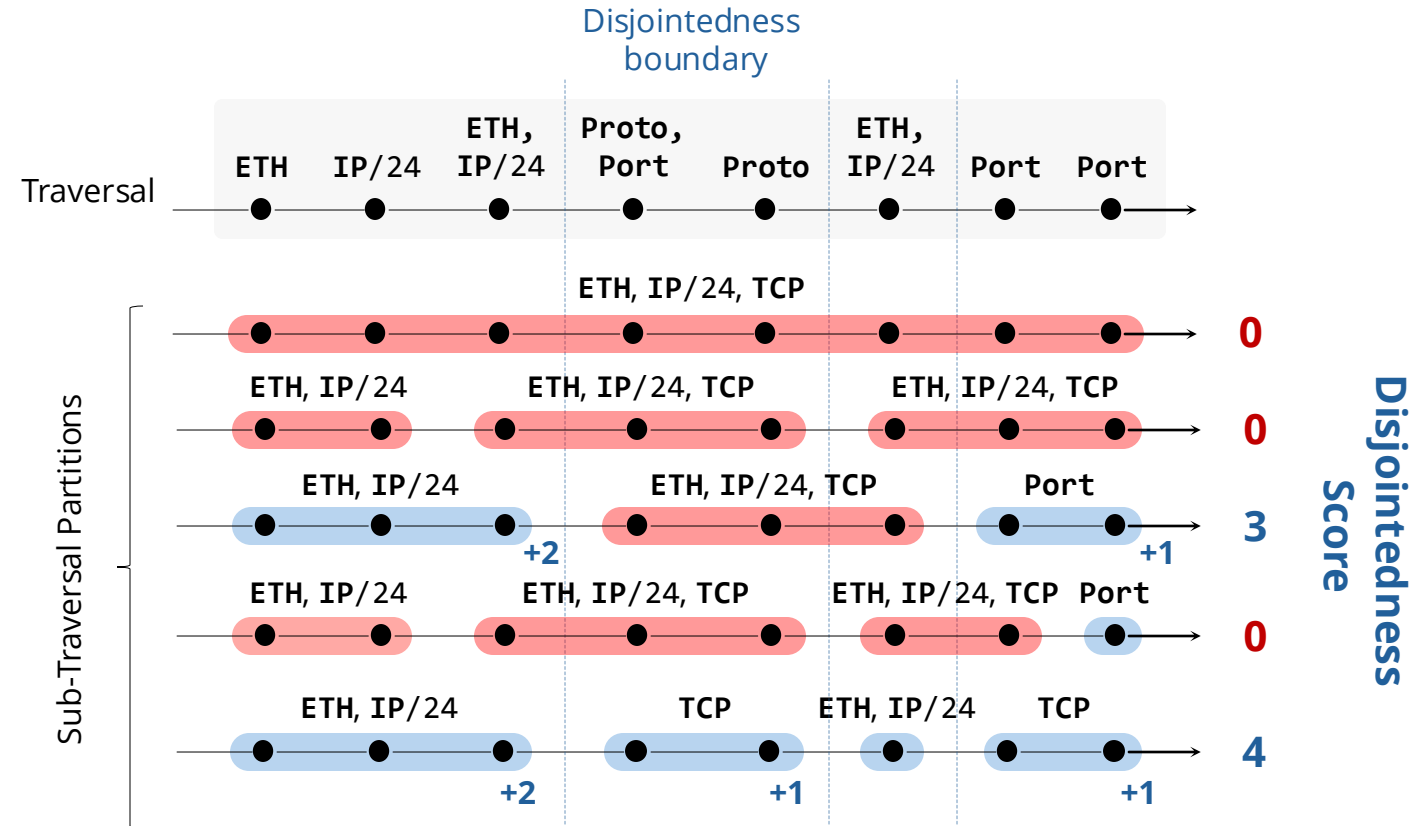
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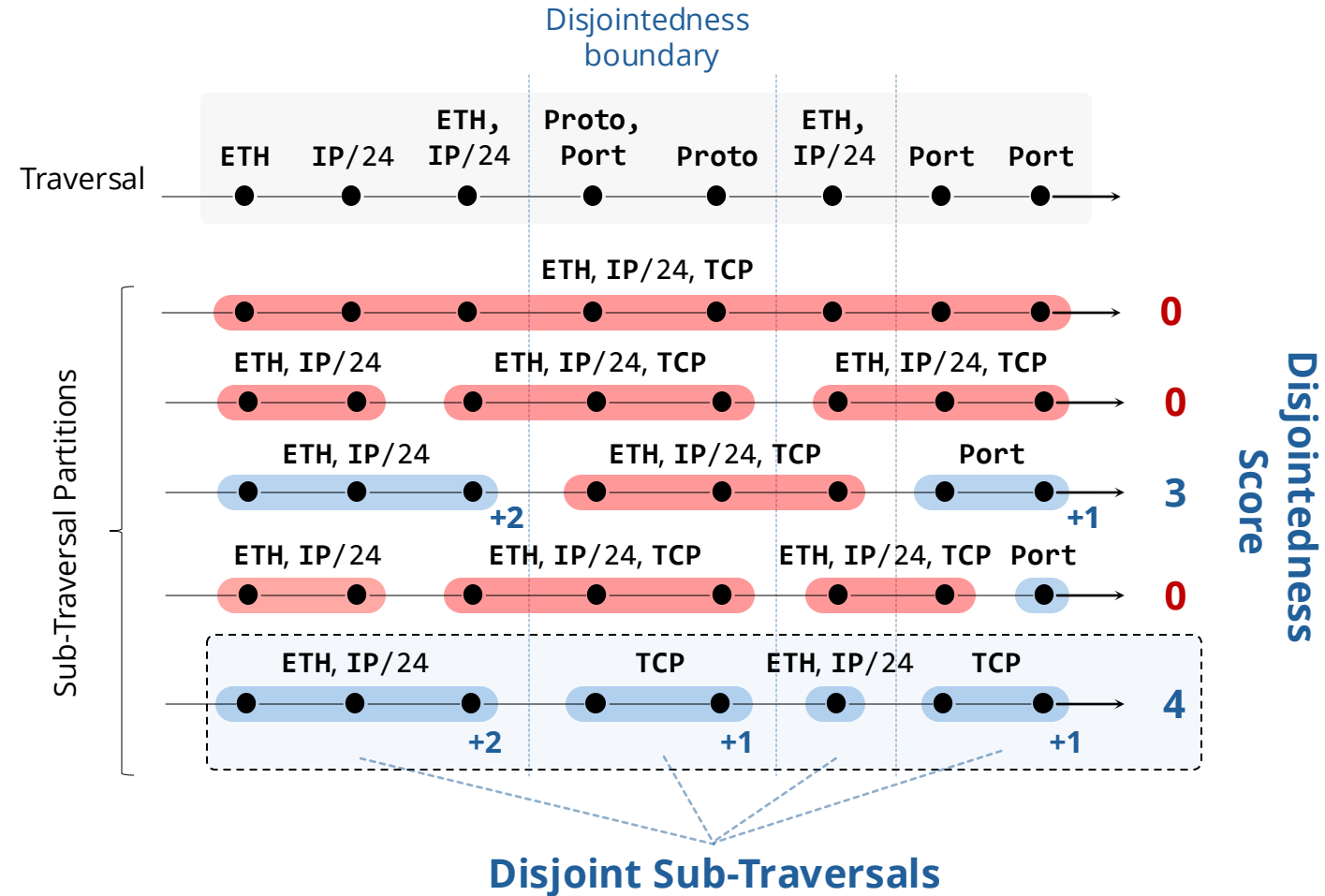
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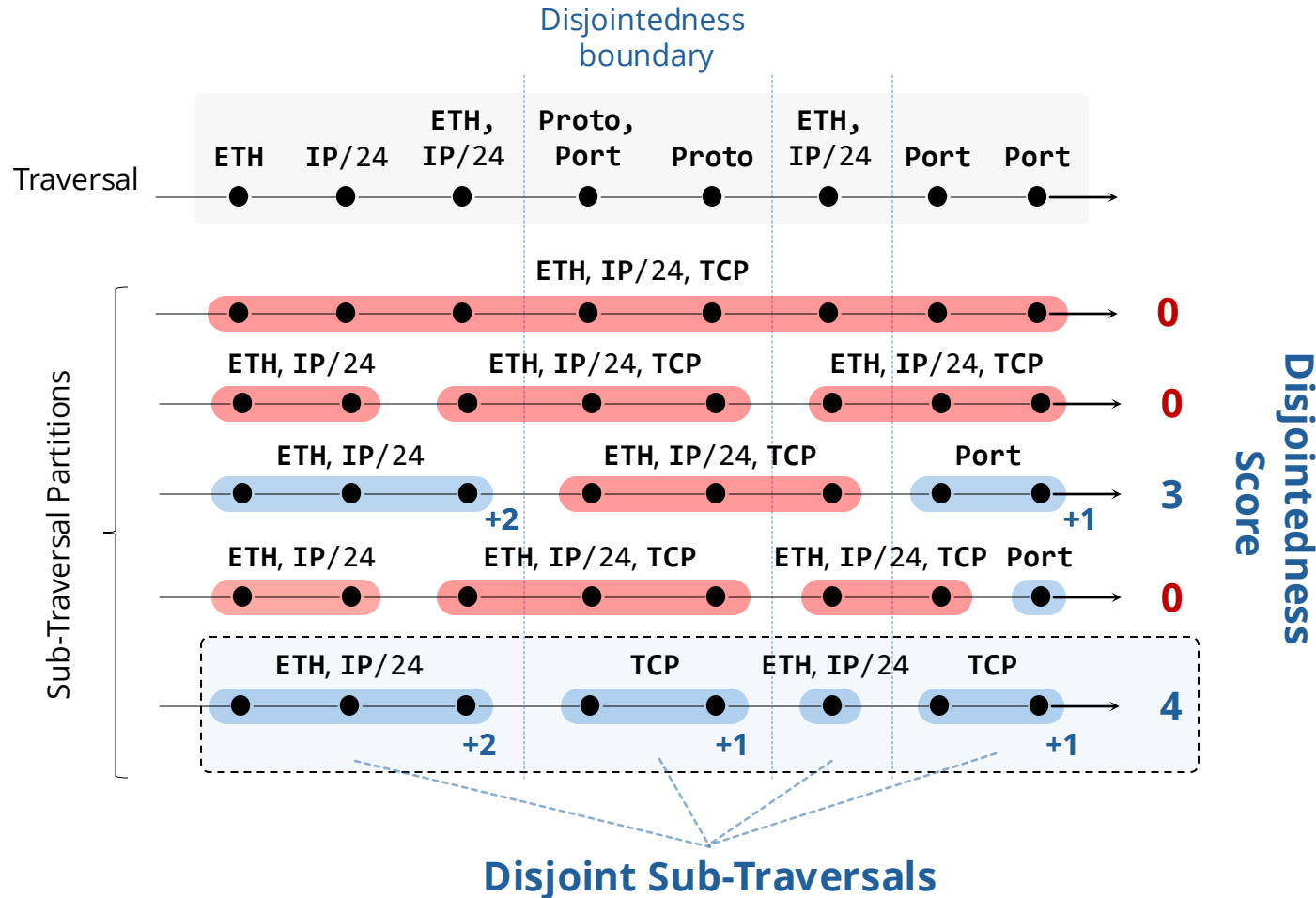
Disjointedness for Pipeline-Aware Locality



Disjointedness for Pipeline-Aware Locality



Disjointedness for Pipeline-Aware Locality



Worst Case Overhead

5x

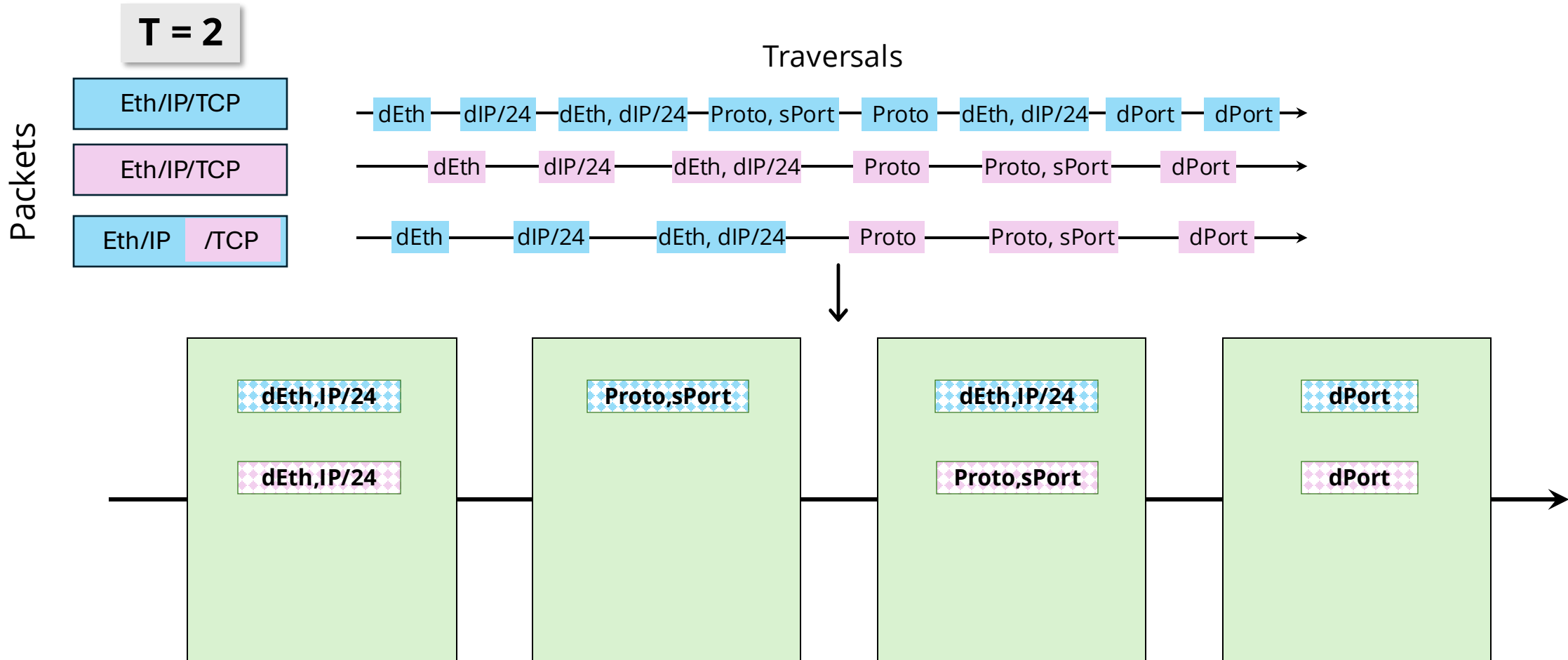
Amortized Overhead

2x

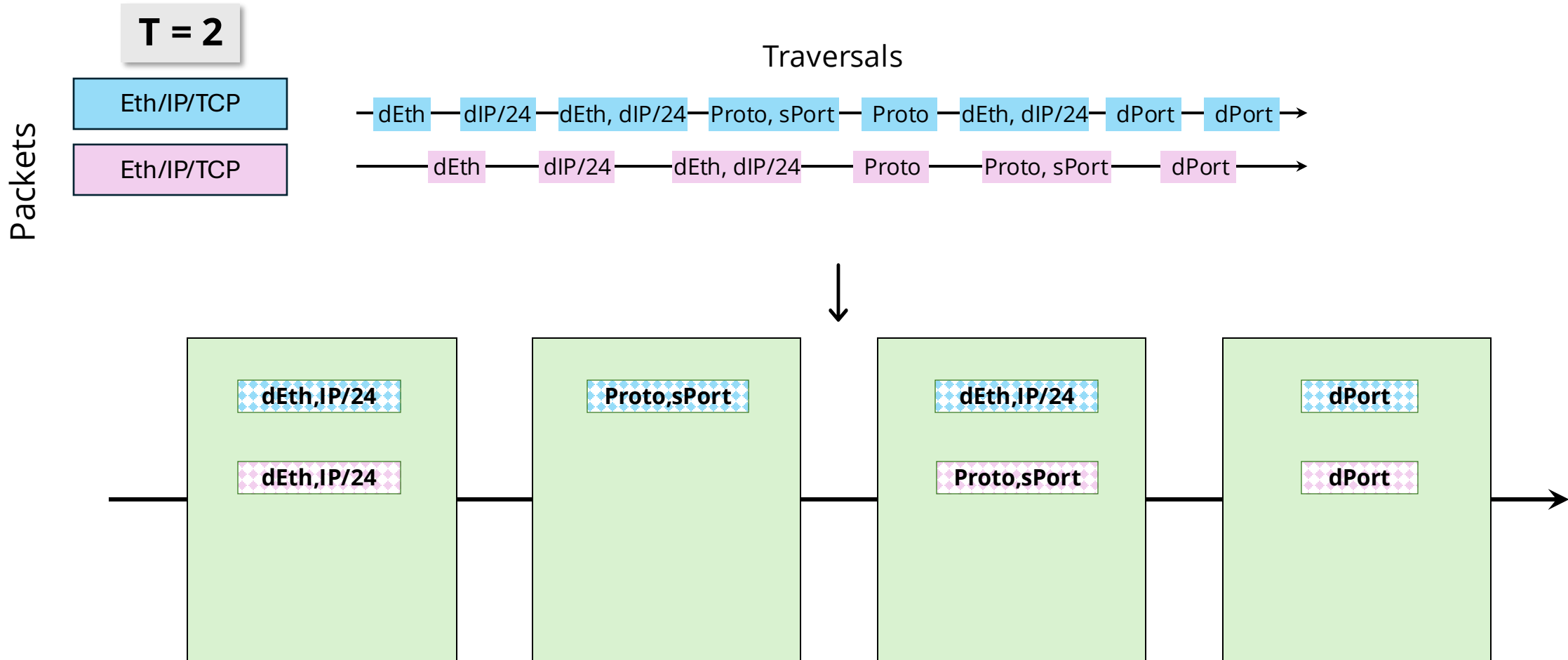
Software processing overhead of Disjoint Partitioning in Gigaflow

Overall cost is negligible because 90% of cache misses are reduced

Performing a Cache Lookup with Gigaflow



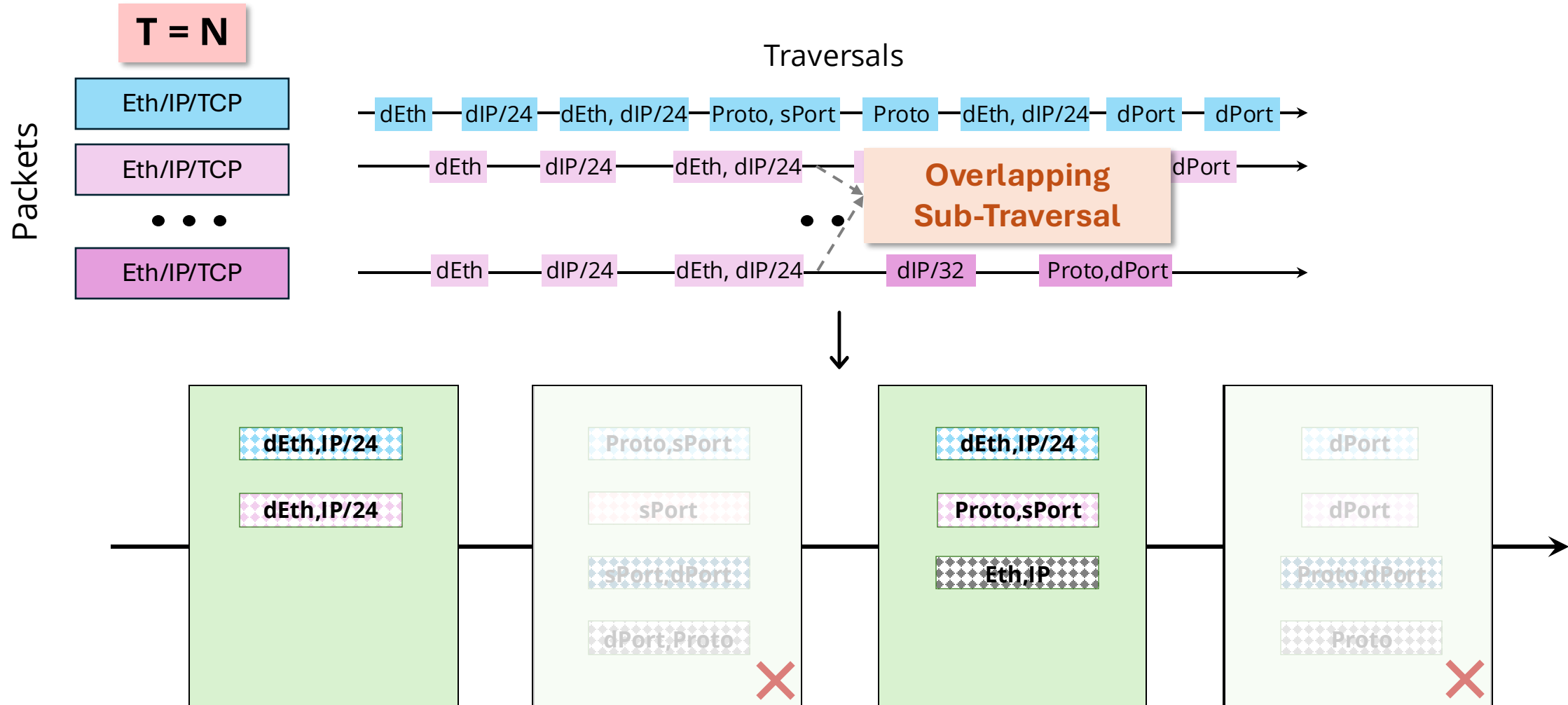
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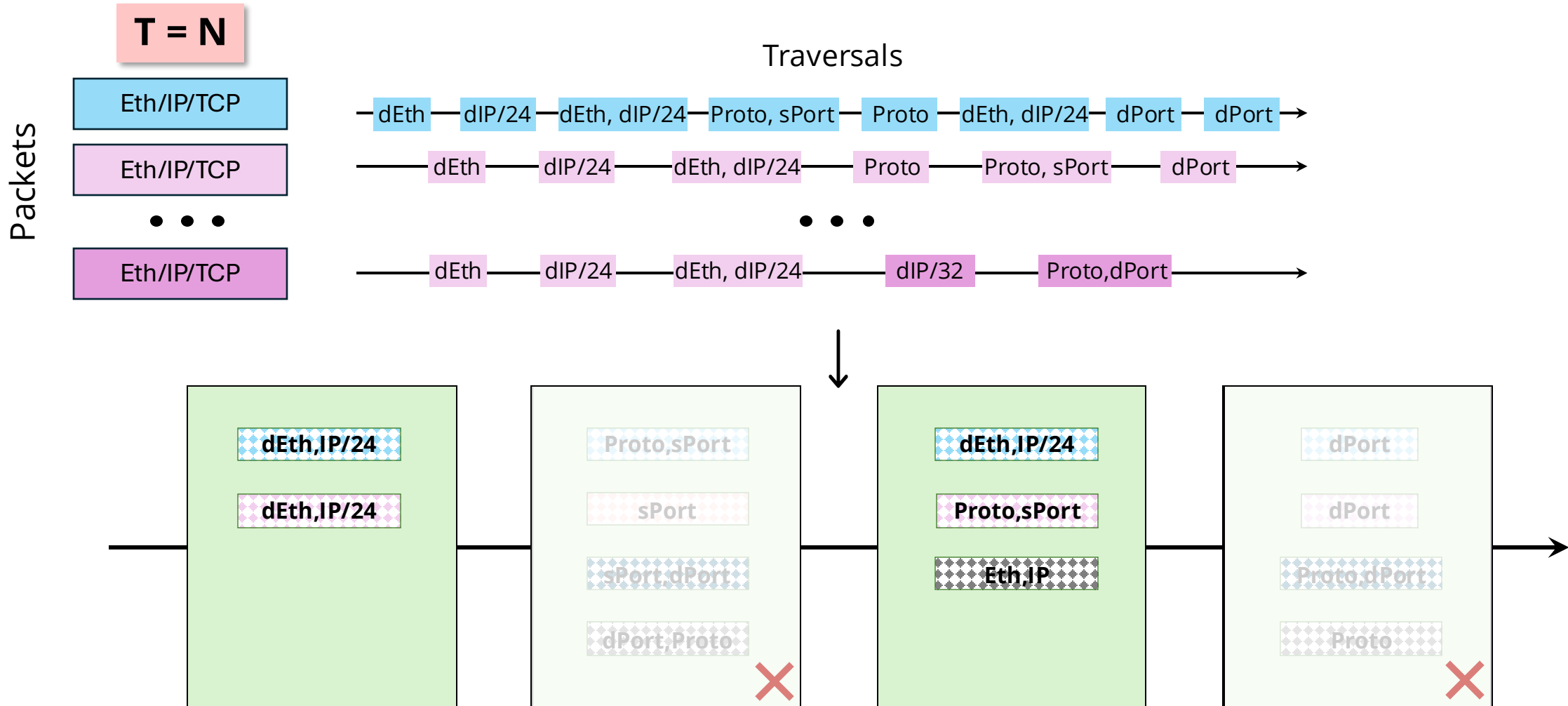
T = N



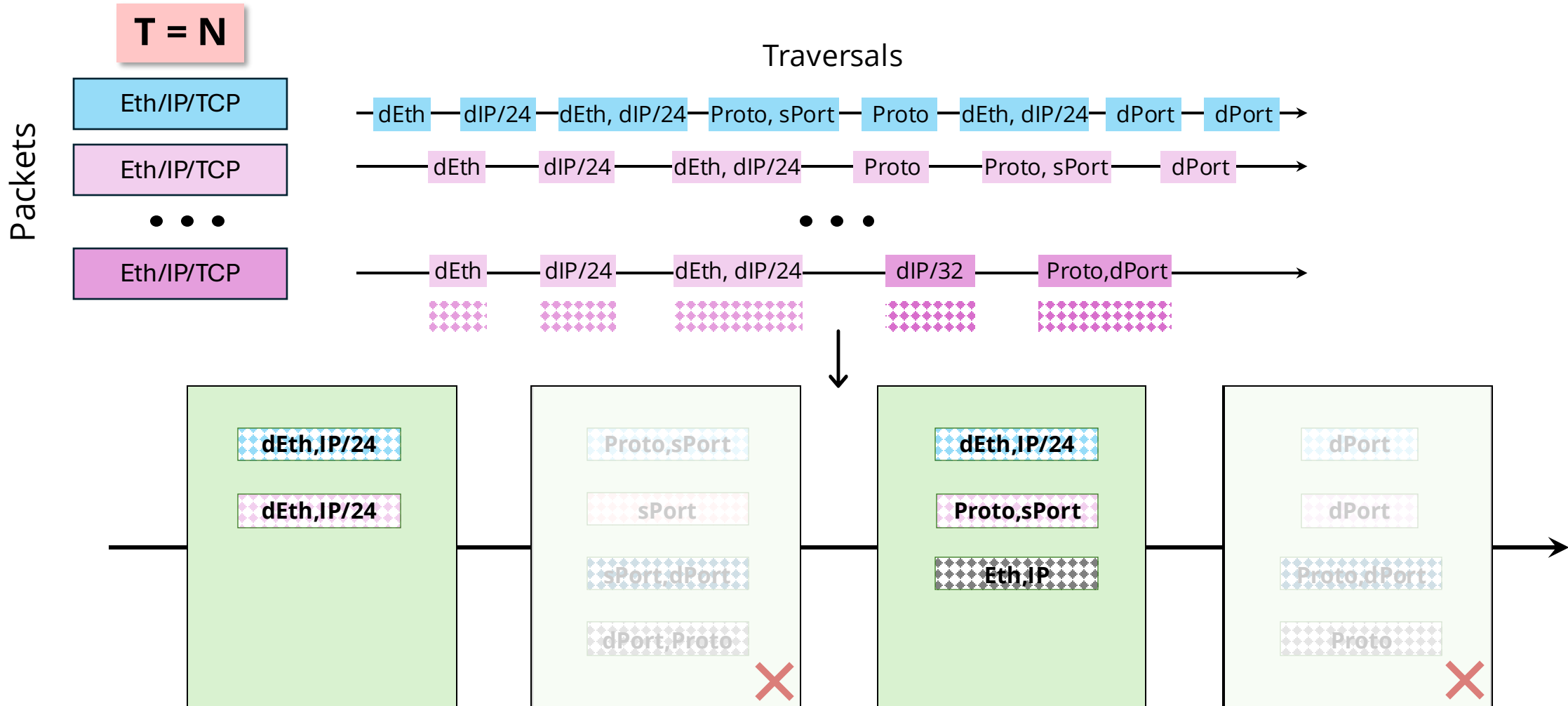
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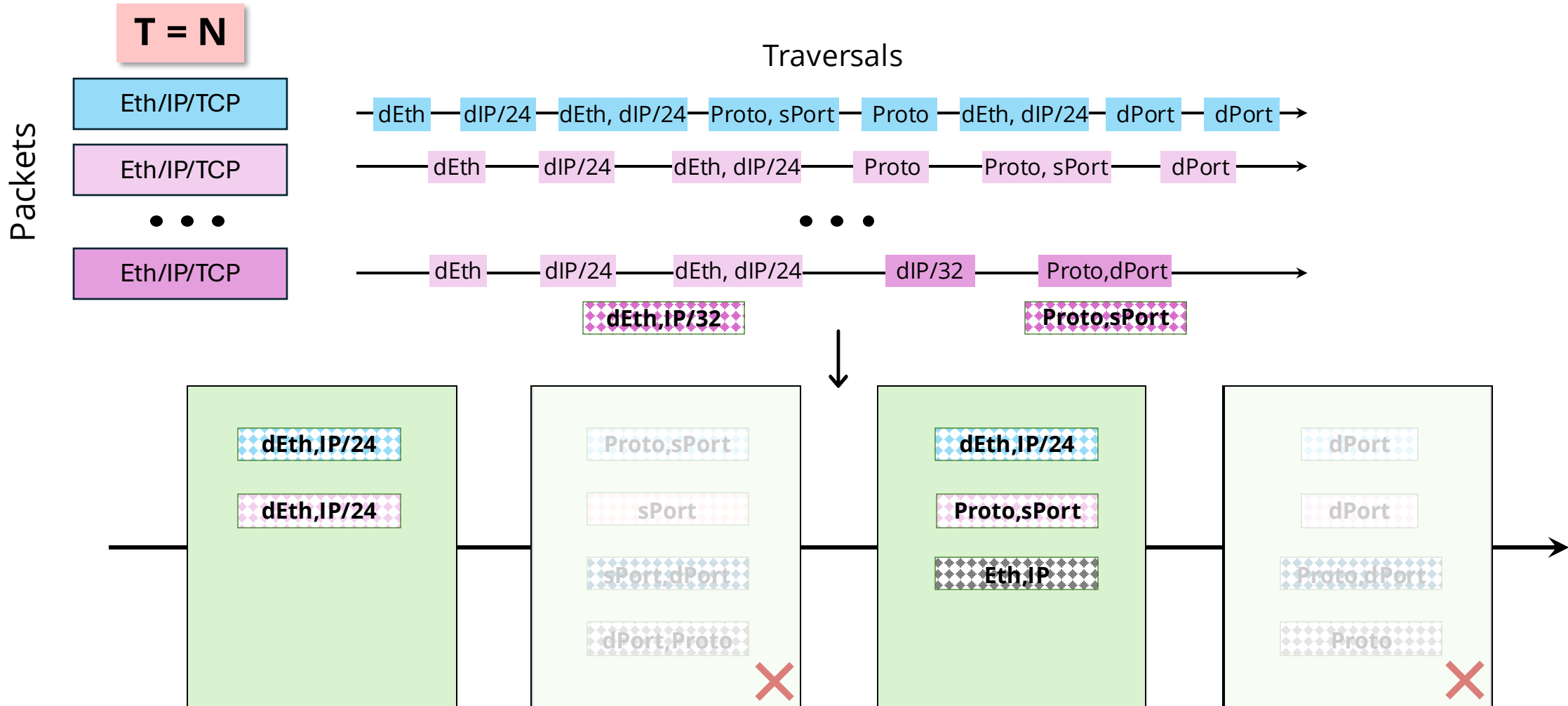
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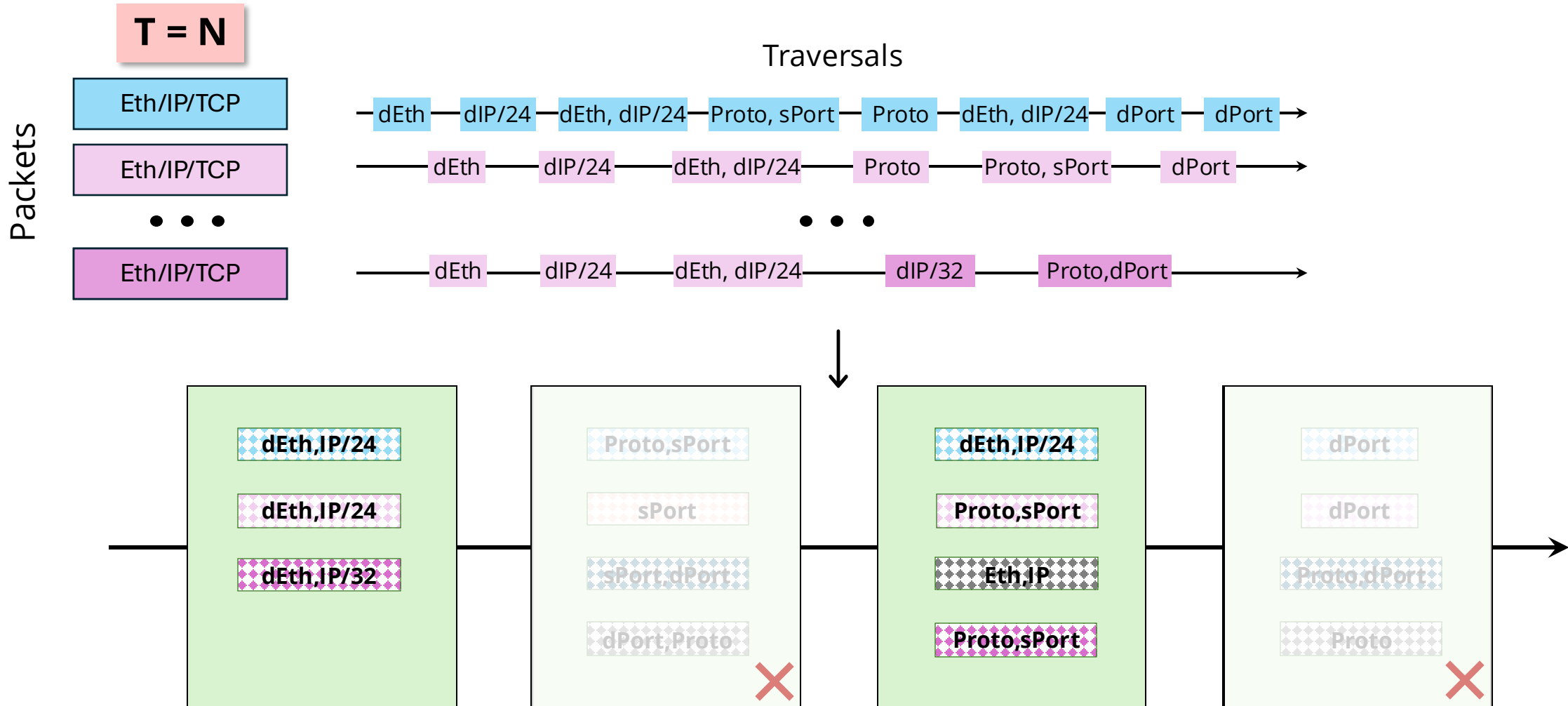
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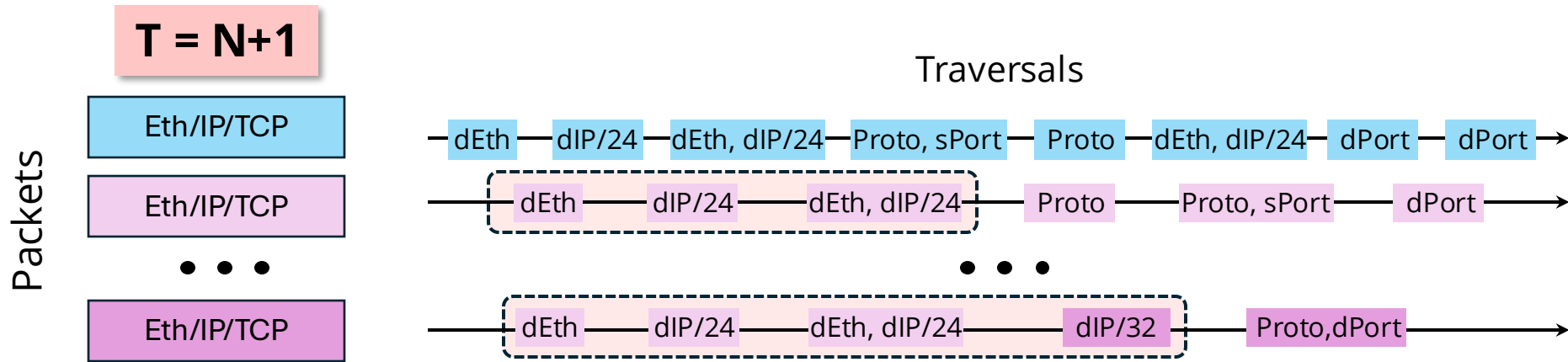
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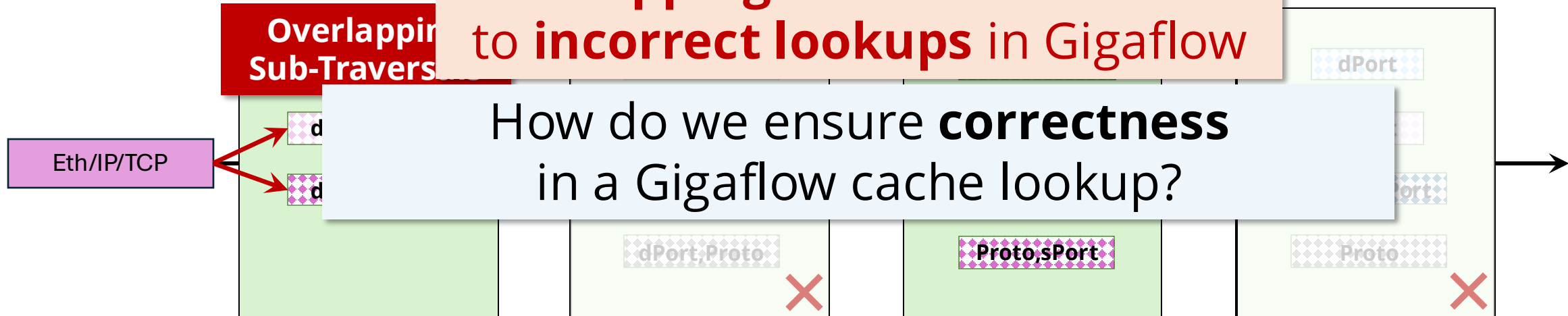


Performing a Cache Lookup with Gigaflow



Overlapping sub-traversals lead to incorrect lookups in Gigaflow

How do we ensure **correctness** in a Gigaflow cache lookup?



Overlapping Rules in Routing Tables

Longest Prefix Match (LPM) in Routing

- **Challenge:**
 - A packet may match multiple destination prefixes
- **Goal:**
 - Route to the best-known destination
- **Rationale:**
 - Most specific match represents *targeted intent*
 - General matches are a *fallback* mechanism
- **Result:**
 - The packet follows the most precise path

10.1.2.130

Table Rules

10.1.2.128/25 → **P:2**

10.1.2.0/24 → **P:4**

10.1.0.0/16 → **P:3**

10.0.0.0/8 → **P:5**

Overlapping Sub-Traversals in Gigaflow

- **Challenge:**

- Packets match multiple overlapping sub-traversals

- **Goal:**

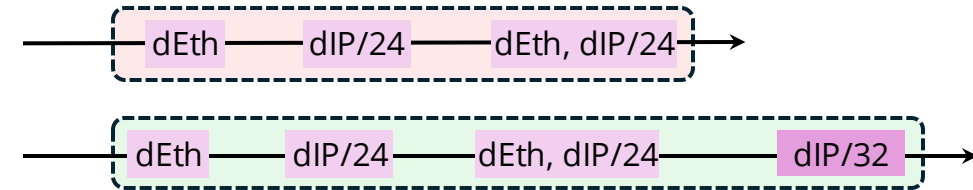
- Highest priority processing for this packet

- **Rationale:**

- Capture processing that might not appear later
- Shorter sub-traversals are a *fallback* mechanism

- **Result:**

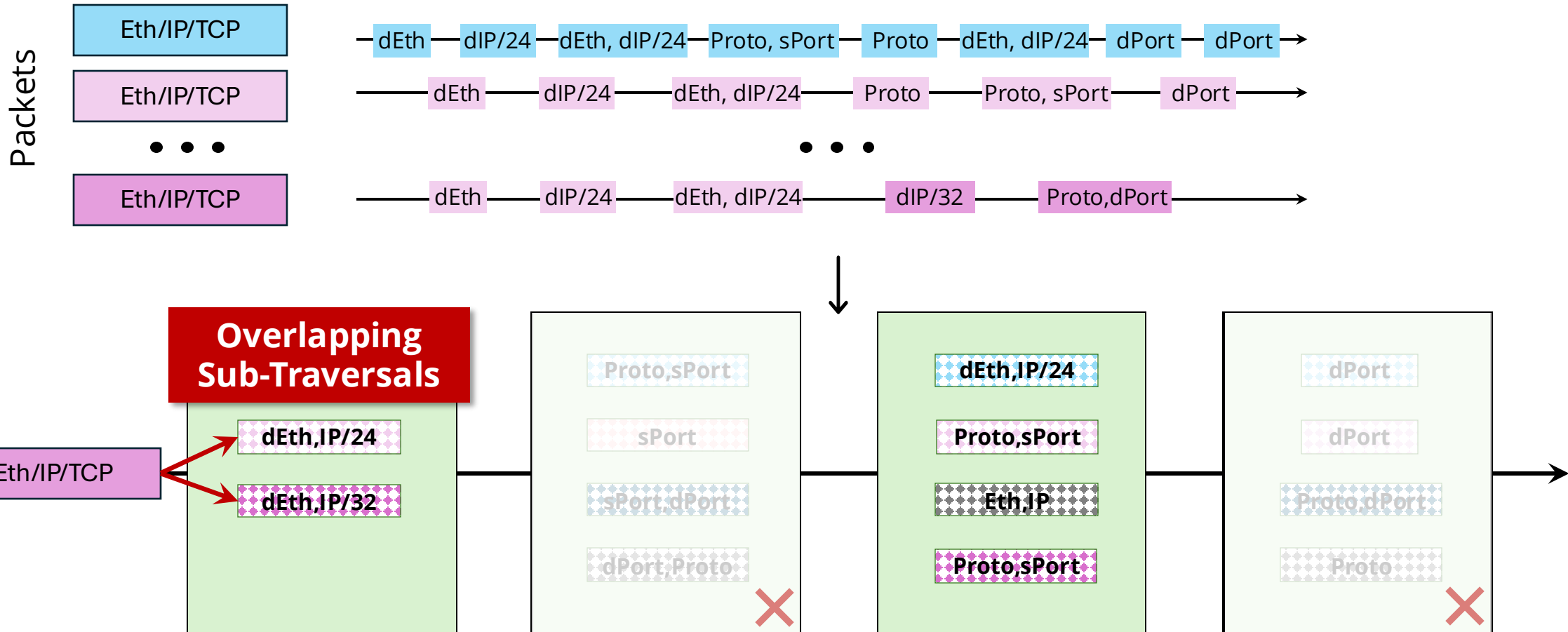
- Gigaflow matches highest priority traversals



Longest Traversal Matching (LTM)

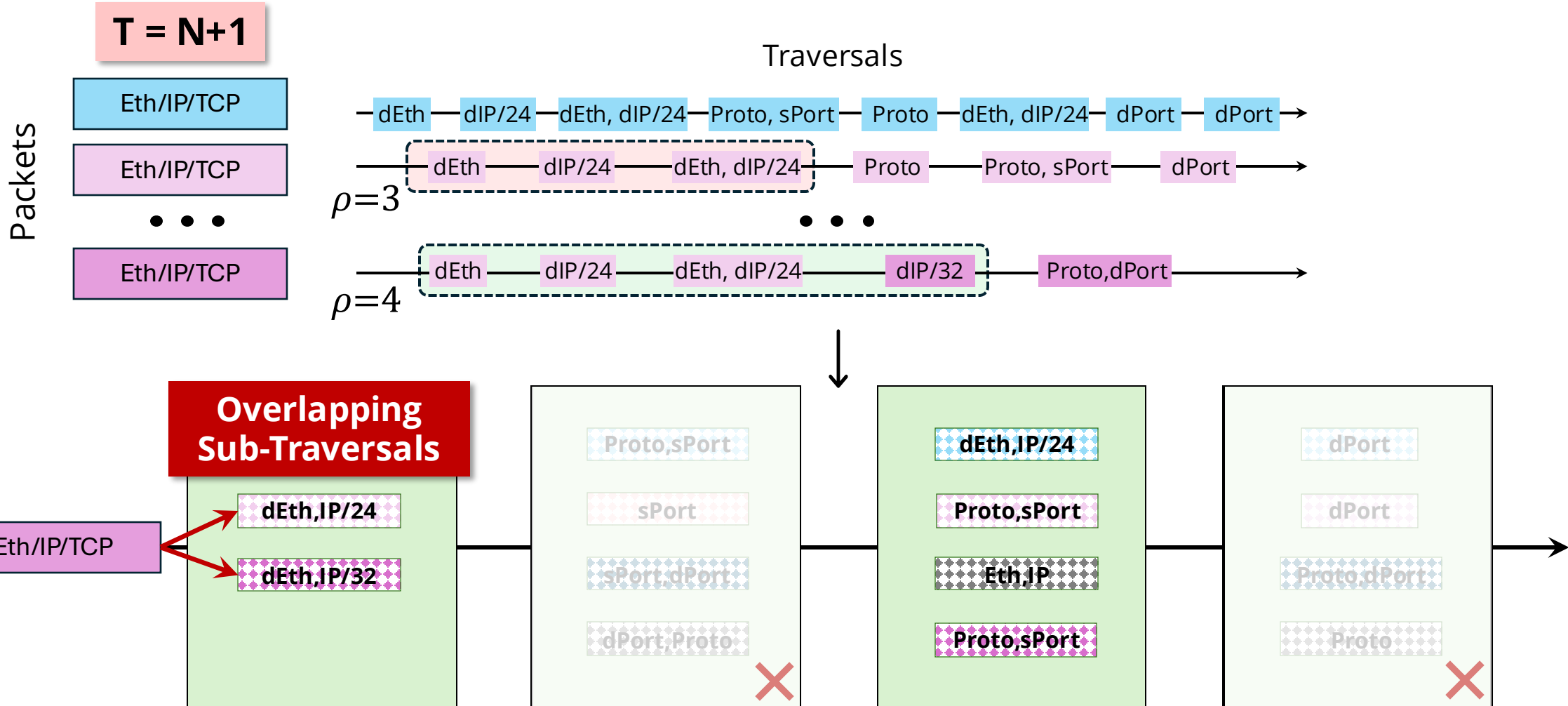
Performing a Cache Lookup with Gigaflow

$$T = N + 1$$



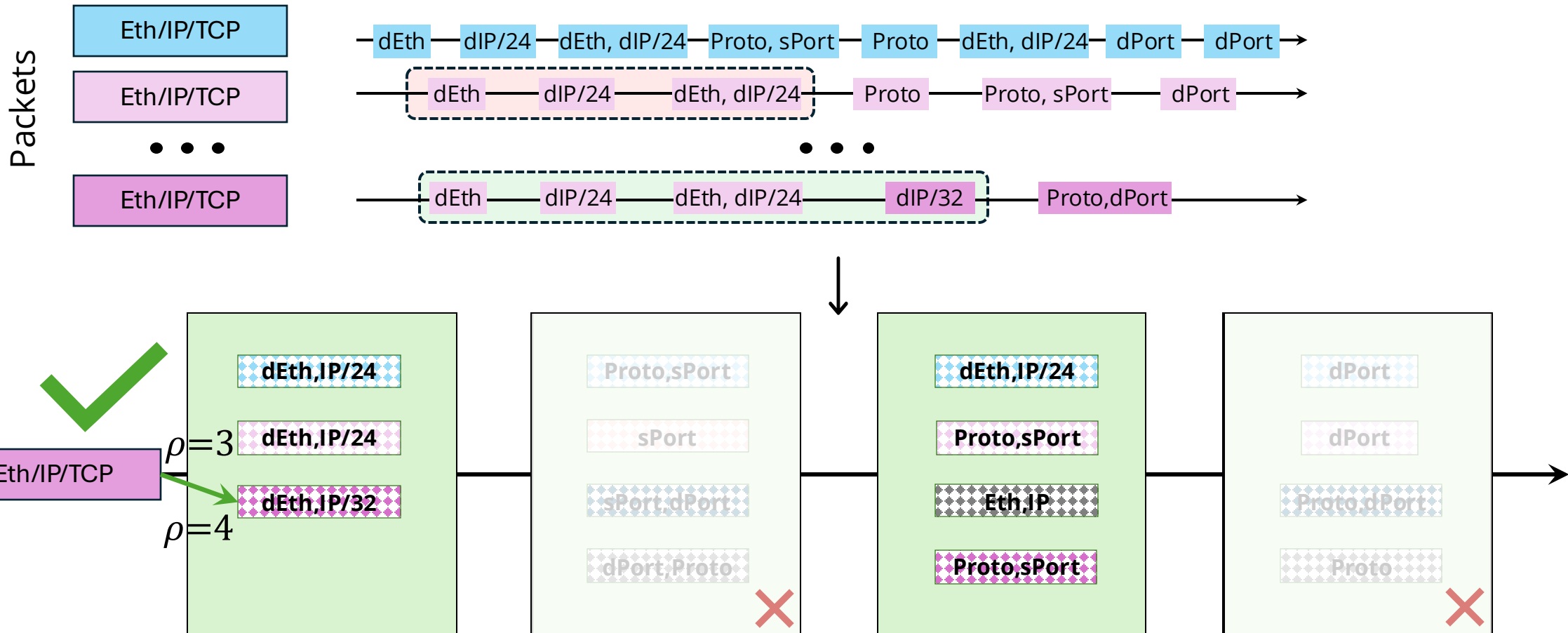
Performing a Cache Lookup with Gigaflow

$$T = N + 1$$



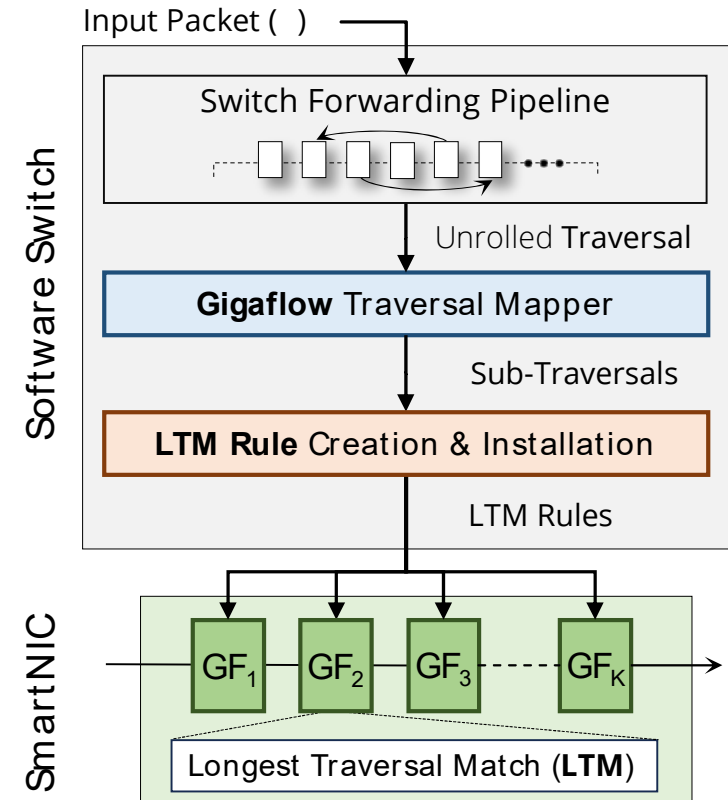
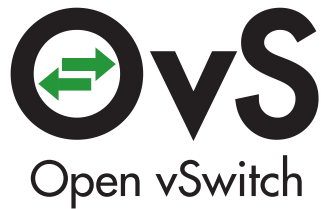
Performing a Cache Lookup with Gigaflow

$$T = N + 1$$



Gigaflow Cache Implementation

- Integrated in the **Open vSwitch**
- Cache offload for **Xilinx Alveo U250**
- **Xilinx P4 SDNet** for hardware tables



High-level view of the Gigaflow in OVS

vSwitch Pipelines and Traffic Workloads

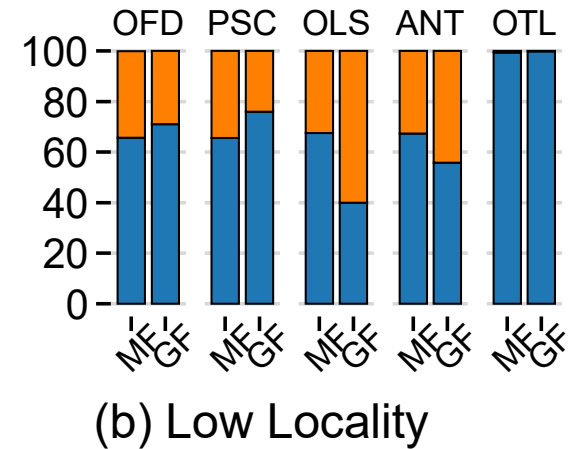
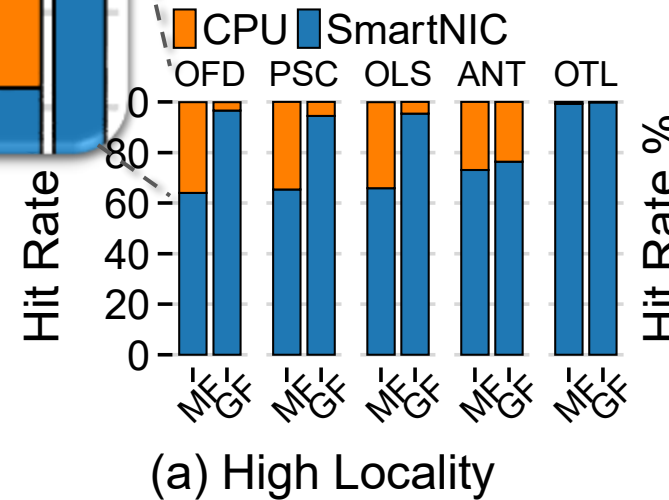
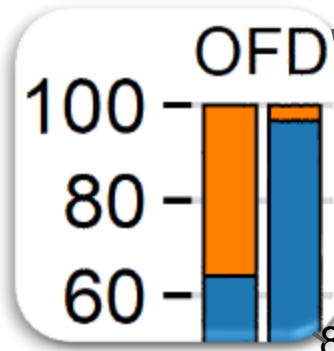
Pipeline	Description	Tables	Traversals
P1: OFD	CORD's Openflow data plane abstraction (OFDPA) for HW/SW switch integration	10	5
P2: PSC	An example L2L3-ACL pipeline implemented for the Pisces paper	7	2
P3: OLS	OVN Logical Switch pipeline to manage logical virtual network topologies in OVS	30	23
P4: ANT	Antrea pipeline implementing networking and security for Kubernetes clusters	22	20
P5: OTL	Openflow Table Type Patterns (TTP) to configure L2L3-ACL policies using OVS	8	11

- Real-world pipeline configurations
 - tables, matching fields, traversals
- Classbench rulesets to populate multi-table rules
- CAIDA traces for realistic traffic patterns
- High/low locality environments
 - high/low sub-traversal sharing opportunity

Performance Summary of Gigaflow

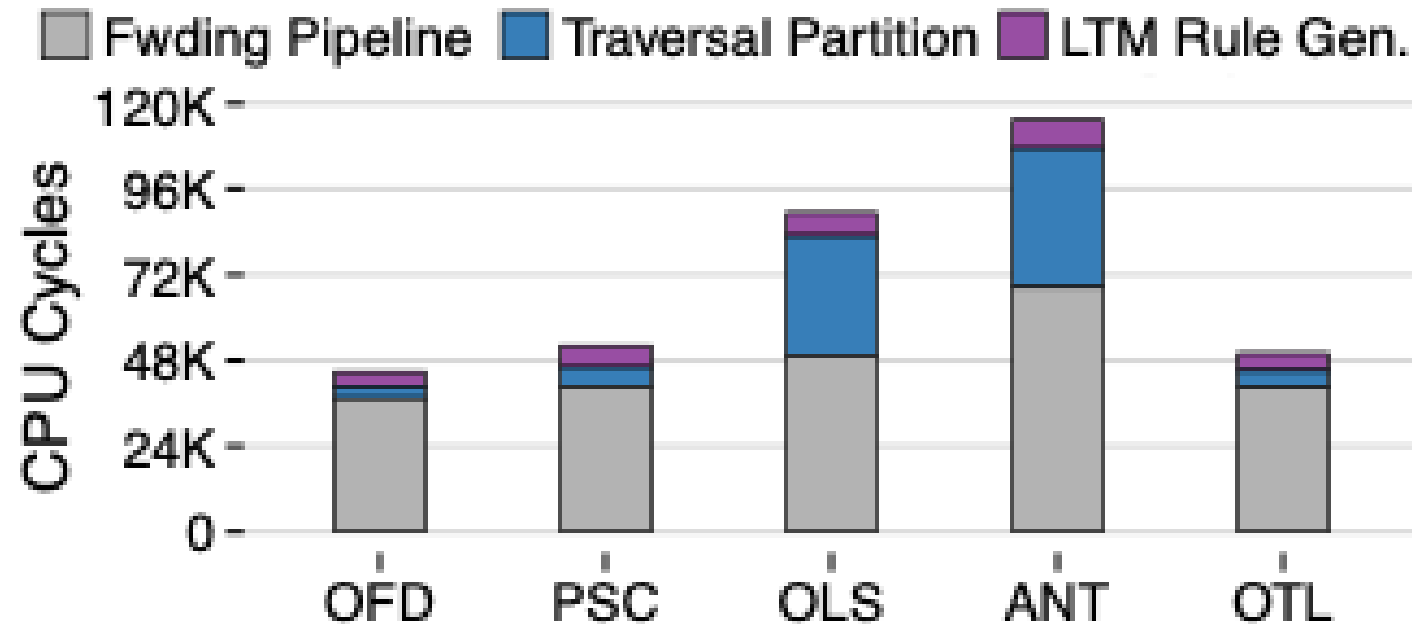
Gigaflow improves

- **#Cache misses** by **90%**
- **Cache hit rate** by **51%**
- **Rule space coverage** by **450x**
- **#Cache entries** by **18%**
- **Forwarding latency** by **30%**



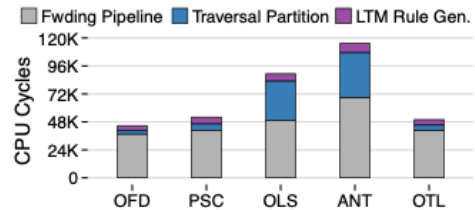
Cache hit rate of Megaflow 32K (MF)
vs Gigaflow 4x8K (GF) for real-world pipelines

Software Processing Overhead of Gigaflow

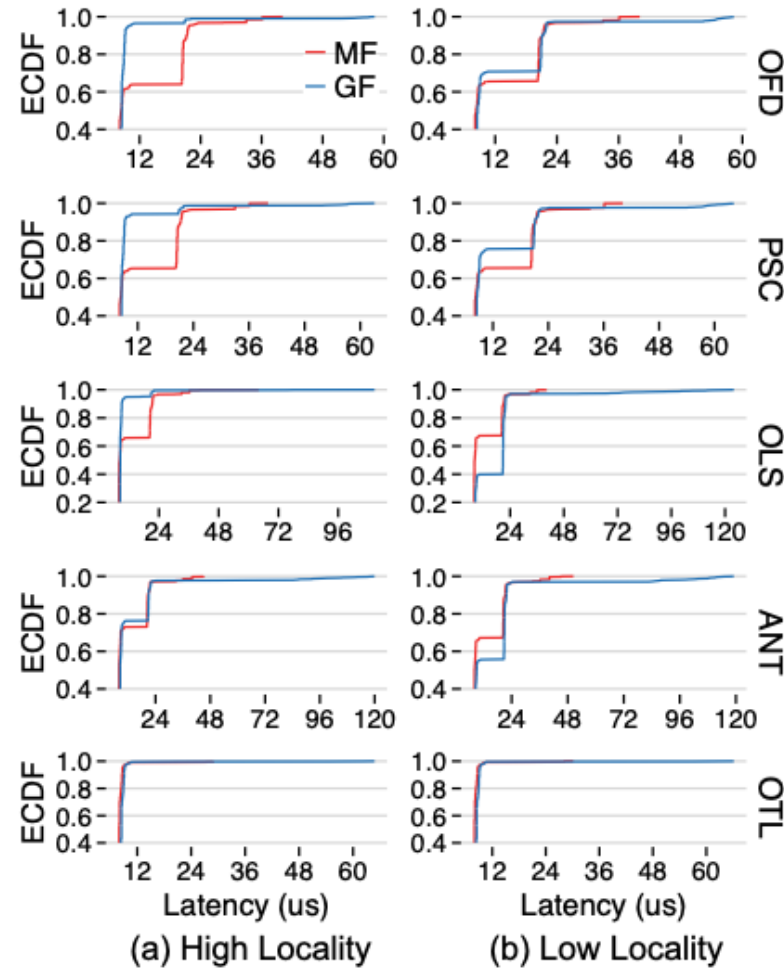


Software overhead of Gigaflow 4x8K (GF)
for five real-world pipelines

End-to-End Packet Latency with Gigaflow

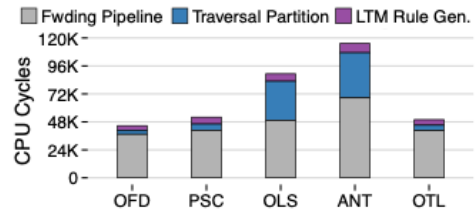


Software overhead of Gigaflow 4x8K (GF) for five real-world pipelines

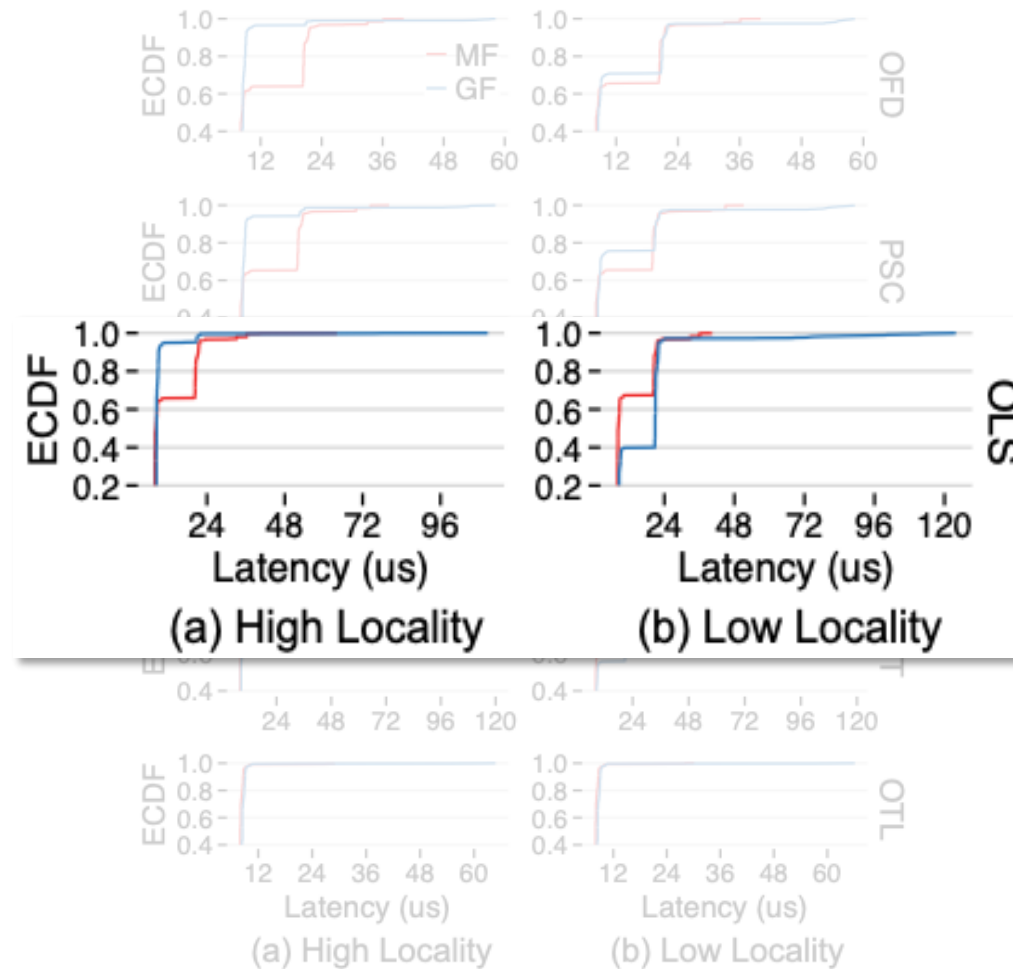


End-to-end latency: Megaflow 32K (MF) vs Gigaflow 4x8K (GF) in high/low locality

End-to-End Packet Latency with Gigaflow

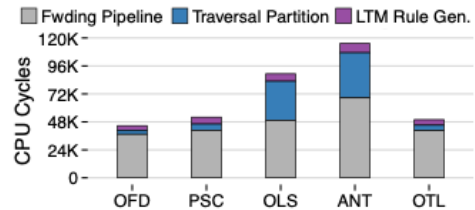


Software overhead of Gigaflow 4x8K (GF) for five real-world pipelines

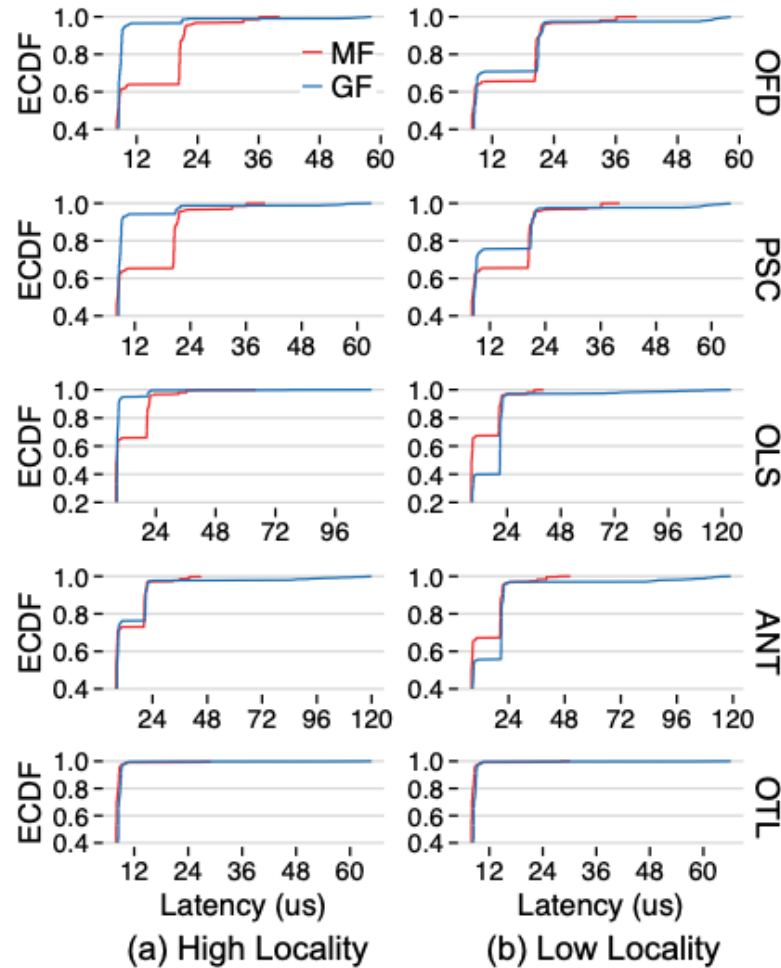


End-to-end latency: Megaflow 32K (MF) vs Gigaflow 4x8K (GF) in high/low locality

End-to-End Packet Latency with Gigaflow

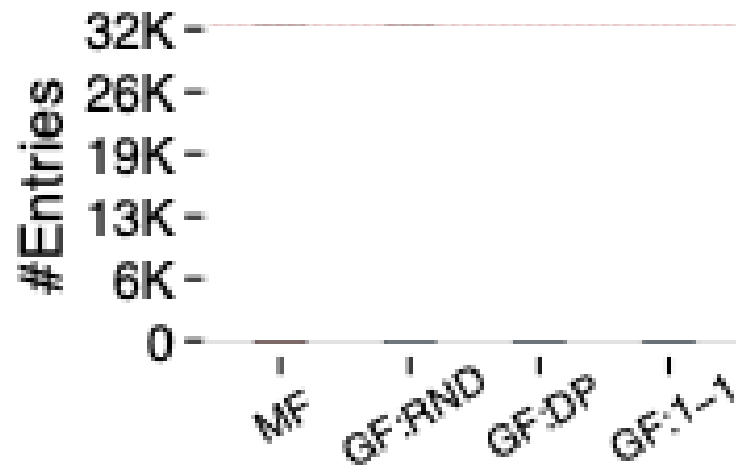


Software overhead of Gigaflow 4x8K (GF) for five real-world pipelines

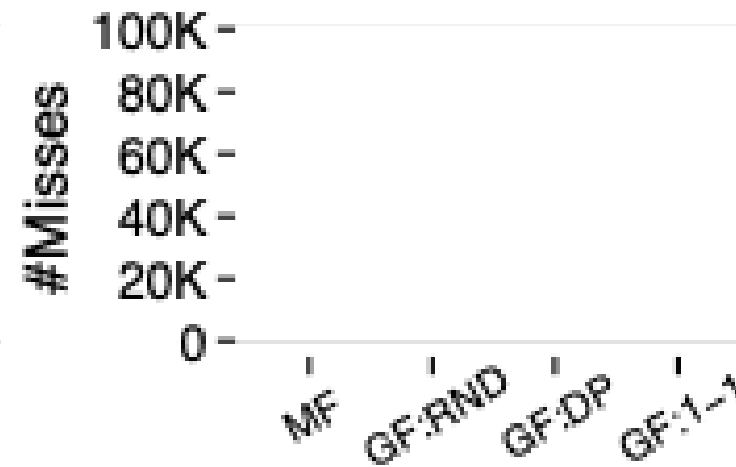


End-to-end latency: Megaflow 32K (MF) vs Gigaflow 4x8K (GF) in high/low locality

Gigaflow with Various Partitioning Strategies



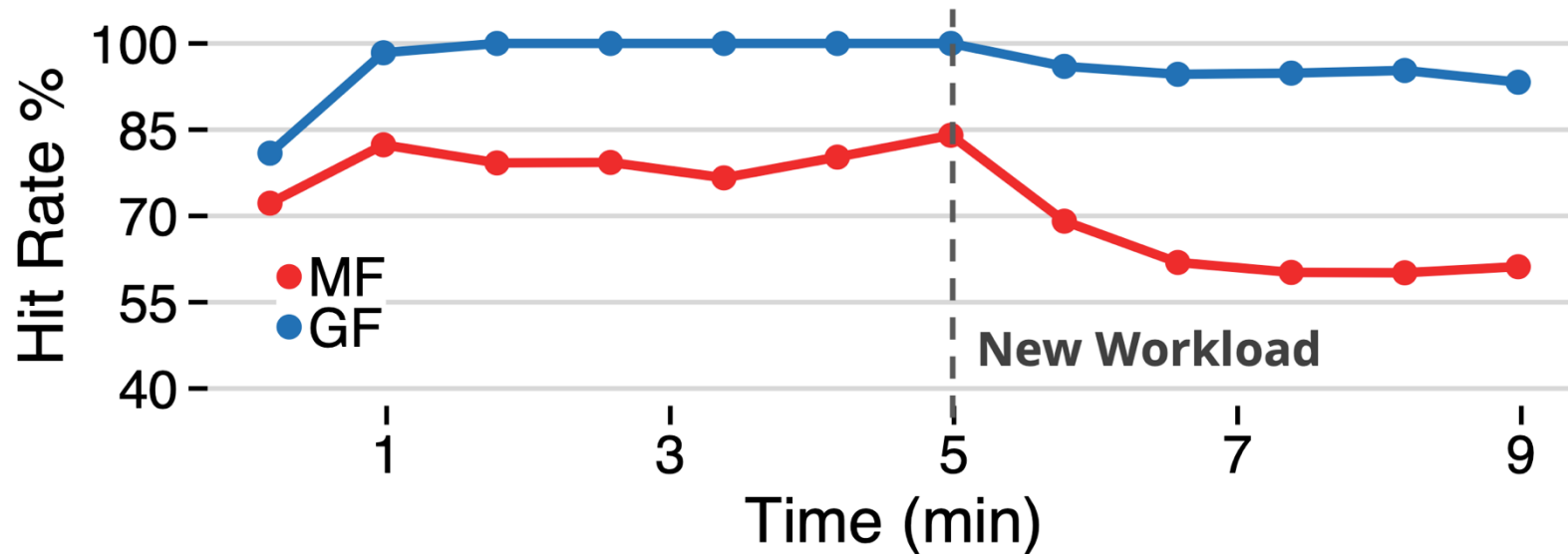
(a) No. of Cache Entries



(b) No. of Misses

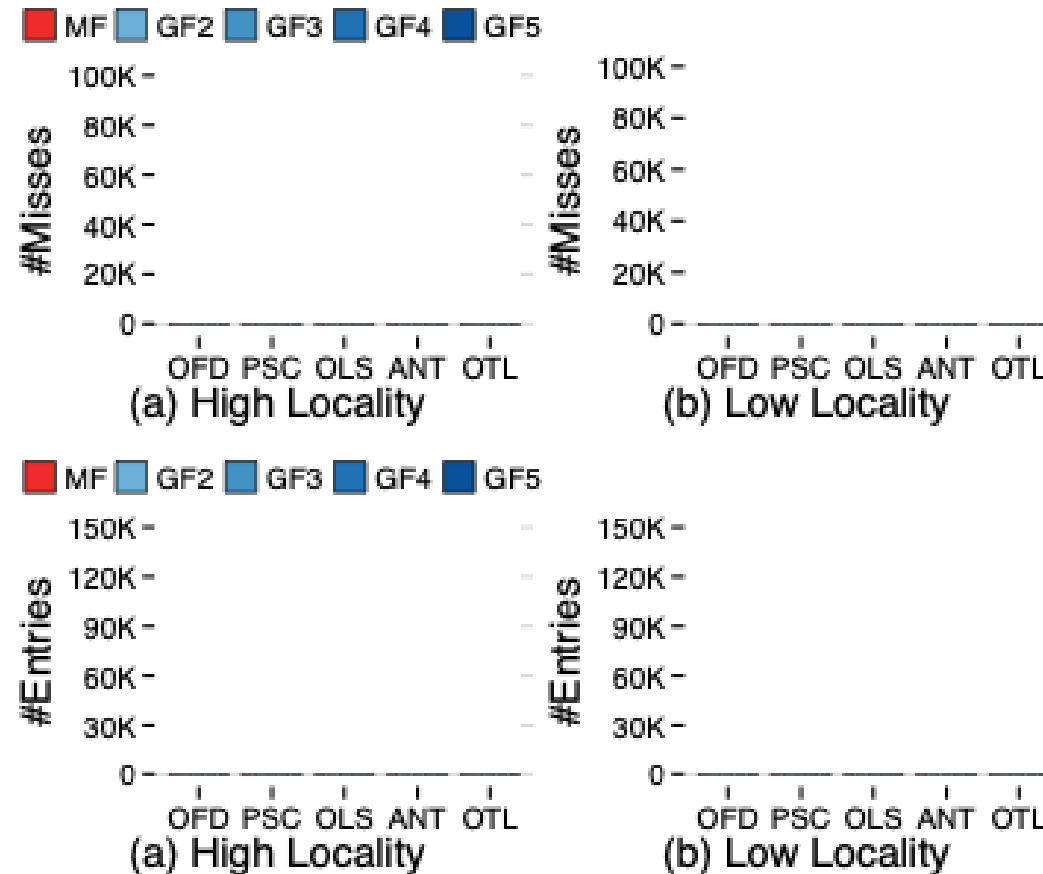
Number of cache entries and cache misses with various partitioning strategies in Gigaflow (GF)

Gigaflow With Dynamic Workloads



Cache hit rate of Megaflow 32K (MF)
vs Gigaflow 4x8K (GF) under dynamic workloads

Gigaflow with Increasing Tables



Cache misses and cache entries with increasing #of Gigaflow (GF) tables

Thank you!

- Gigaflow → a *sub-traversal* cache
- Captures **pipeline-aware** locality
- Delivers **higher rule space coverage**



NextGArch Lab – <https://nextgarch.github.io/>

Contact us: zulfiqaa@umich.edu

